
Playing Infinite Games in Finite Time

Joint work with John Fearnley,
Daniel Neider, and Roman Rabinovich

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RWTH Aachen University

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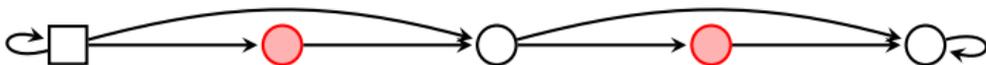
Introduction

How can you play infinite games in finite time? Aren't they infinite for a reason?

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Introductory example

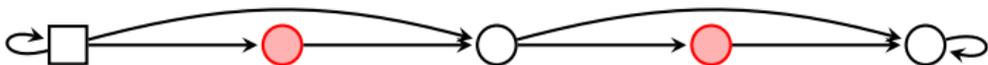


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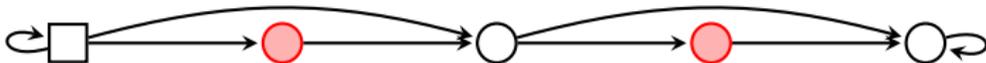
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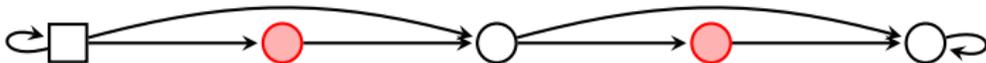
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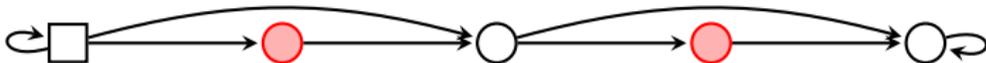
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Can we play in finite time without relying on a memory structure?

Muller Games

Inspired by previous work of McNaughton on playing infinite games in finite time, we consider Muller games $(\mathcal{A}, \mathcal{F}_0, \mathcal{F}_1)$:

- arena \mathcal{A} and partition $(\mathcal{F}_0, \mathcal{F}_1)$ containing the loops of \mathcal{A} .
- Player i wins ρ iff $\text{Inf}(\rho) = \{v \mid \exists^\omega n \text{ s.t. } \rho_n = v\} \in \mathcal{F}_i$.

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Running example



- $\mathcal{F}_0 = \{\{0, 1, 2\}, \{0\}, \{2\}\}$
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Player 0 has a winning strategy from every vertex: alternate between 0 and 2.

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Remark

Muller games are not reducible to safety games.

Outline

1. **Playing Muller Games in Finite Time**
2. Solving Muller Games by Solving Safety Games
3. Conclusion

Scoring Functions

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$$\text{Sc}_F(v) = \begin{cases} 1 & \text{if } F = \{v\}, \\ 0 & \text{otherwise,} \end{cases}$$

and

$$\text{Acc}_F(v) = \begin{cases} \emptyset & \text{if } F = \{v\}, \\ F \cap \{v\} & \text{otherwise.} \end{cases}$$

Scoring Functions

Let $F \subseteq V$, $F \neq \emptyset$. For $v \in V$ and $w \in V^+$ define

$$\text{Sc}_F(wv) = \begin{cases} 0 & \text{if } v \notin F, \\ \text{Sc}_F(w) & \text{if } v \in F \wedge \text{Acc}_F(w) \neq (F \setminus \{v\}), \\ \text{Sc}_F(w) + 1 & \text{if } v \in F \wedge \text{Acc}_F(w) = (F \setminus \{v\}), \end{cases}$$

and

$$\text{Acc}_F(wv) = \begin{cases} \emptyset & \text{if } v \notin F, \\ \text{Acc}_F(w) \cup \{v\} & \text{if } v \in F \wedge \text{Acc}_F(w) \neq (F \setminus \{v\}), \\ \emptyset & \text{if } v \in F \wedge \text{Acc}_F(w) = (F \setminus \{v\}). \end{cases}$$

Scoring Functions cont'd

- $Sc_F(w)$: maximal $k \in \mathbb{N}$ such that F is visited k times since last vertex in $V \setminus F$ (reset).
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w	0	0	1	1	0	0	1	2
$Sc_{\{0,1\}}$	0	0	1	1	2	2	3	0
$Acc_{\{0,1\}}$	{0}	{0}	\emptyset	{1}	\emptyset	{0}	\emptyset	\emptyset
$Sc_{\{0,1,2\}}$	0	0	0	0	0	0	0	1
$Acc_{\{0,1,2\}}$	{0}	{0}	{0, 1}	{0, 1}	{0, 1}	{0, 1}	{0, 1}	\emptyset

Scoring Functions cont'd

- $Sc_F(w)$: maximal $k \in \mathbb{N}$ such that F is visited k times since last vertex in $V \setminus F$ (reset).
- $Acc_F(w)$: set $A \subset F$ of vertices seen since last increase or reset of Sc_F .

Example:

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$Sc_{\{0,1\}}$	0	0	1	1	2	2	3	0
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$Acc_{\{0,1,2\}}$	{0}	{0}	{0, 1}	{0, 1}	{0, 1}	{0, 1}	{0, 1}	\emptyset

Remark

$$F = \text{Inf}(\rho) \Leftrightarrow \liminf_{n \rightarrow \infty} Sc_F(\rho_0 \cdots \rho_n) = \infty$$

Finite-time Muller Games

Two properties of scoring functions (informal versions):

1. If you play long enough (i.e., $k^{|V|}$ steps), some score value will be high (i.e., k).
2. At most one score value can increase at a time.

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Rules:

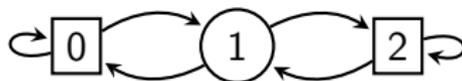
- Players move a token through the arena.
- Stop play w as soon as score of k is reached for the first time.
- There is a unique F such that $\text{Sc}_F(w) = k$ (see above).
- Player i wins w iff $F \in \mathcal{F}_i$.

Two Examples



- $\mathcal{F}_0 = \{\{0, 1, 2\}, \{0\}, \{2\}\}$
- $\mathcal{F}_1 = \{\{0, 1\}, \{1, 2\}\}$

Two Examples

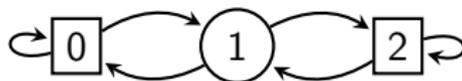


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Losing player (Player 1) can enforce score of two:

Two Examples



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Losing player (Player 1) can enforce score of two:

$$1 \rightarrow 2 \quad (\text{w.l.o.g.})$$

Two Examples



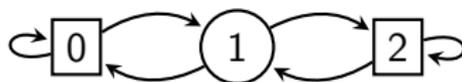
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Losing player (Player 1) can enforce score of two:

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Two Examples

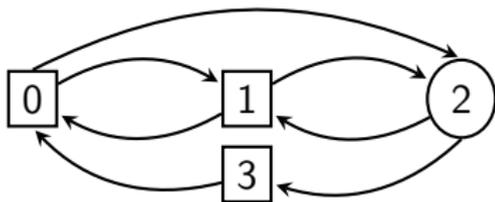


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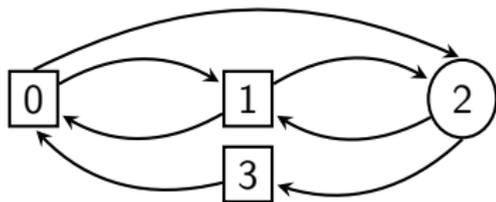


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Losing player (Player 1) is the first to reach a score of two:

Two Examples

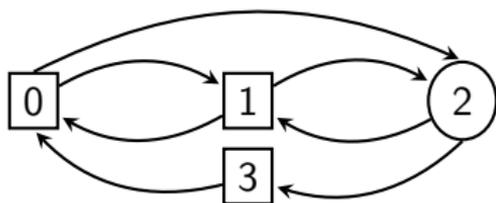


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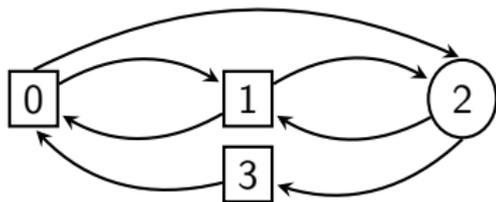


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Losing player (Player 1) is the first to reach a score of two:

$$3 \rightarrow 0$$

Two Examples

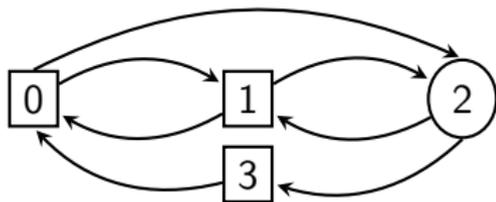


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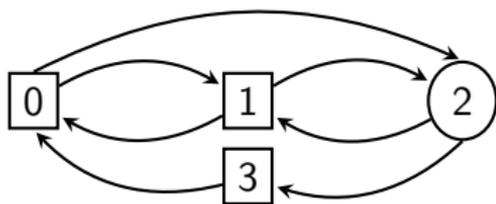


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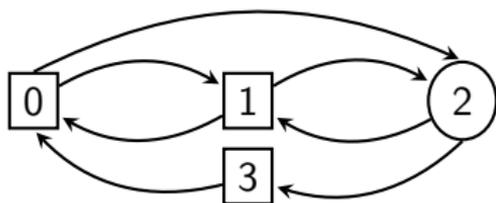


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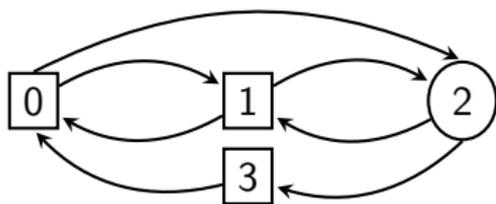


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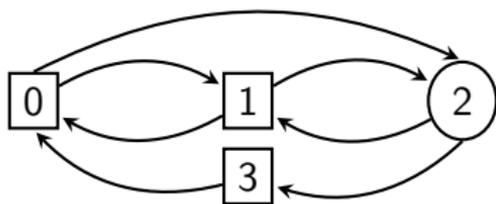


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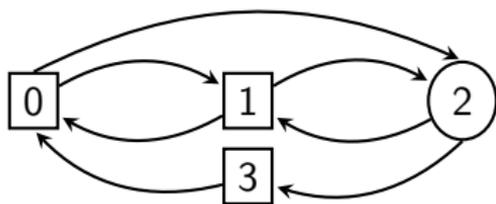


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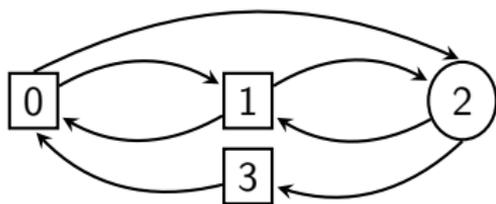


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Losing player (Player 1) is the first to reach a score of two:

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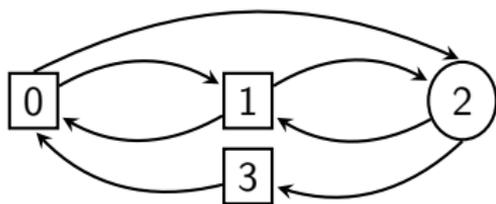


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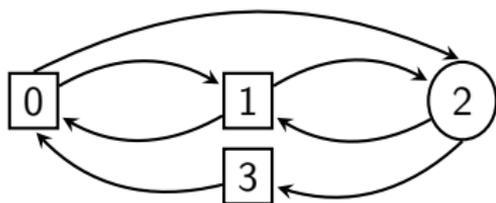


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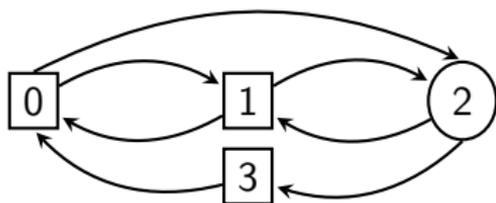


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Theorem (FZ10)

The winning regions in a Muller game $(\mathcal{A}, \mathcal{F}_0, \mathcal{F}_1)$ and in the finite-time Muller game $(\mathcal{A}, \mathcal{F}_0, \mathcal{F}_1, 3)$ coincide.

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Stronger statement, which implies the theorem:

Lemma (FZ10)

On her winning region, Player i can prevent her opponent from ever reaching a score of 3 for every set $F \in \mathcal{F}_{1-i}$.

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Corollary

Two “reductions”: Muller game to ..

- 1. .. reachability game on unravelling up to score 3.*
- 2. .. safety game: see next slides.*

Outline

1. Playing Muller Games in Finite Time
- 2. Solving Muller Games by Solving Safety Games**
3. Conclusion

“Reducing“ Muller games to Safety Games



- $\mathcal{F}_0 = \{\{0, 1, 2\}, \{0\}, \{2\}\}$
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Idea: track of Player 1's scores and avoid $\text{Sc}_F = 3$ for $F \in \mathcal{F}_1$.

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Idea: track of Player 1's scores and avoid $\text{Sc}_F = 3$ for $F \in \mathcal{F}_1$.

- Ignore scores of Player 0.
- Identify plays having the same scores and accumulators for Player 1: $w =_{\mathcal{F}_1} w'$ iff $\text{last}(w) = \text{last}(w')$ and for all $F \in \mathcal{F}_1$:

$$\text{Sc}_F(w) = \text{Sc}_F(w') \text{ and } \text{Acc}_F(w) = \text{Acc}(w')$$

- Build $=_{\mathcal{F}_1}$ -quotient of unravelling up to score 3 for Player 1.
- Winning condition for Player 0: avoid $\text{Sc}_F = 3$ for all $F \in \mathcal{F}_1$.

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[0]

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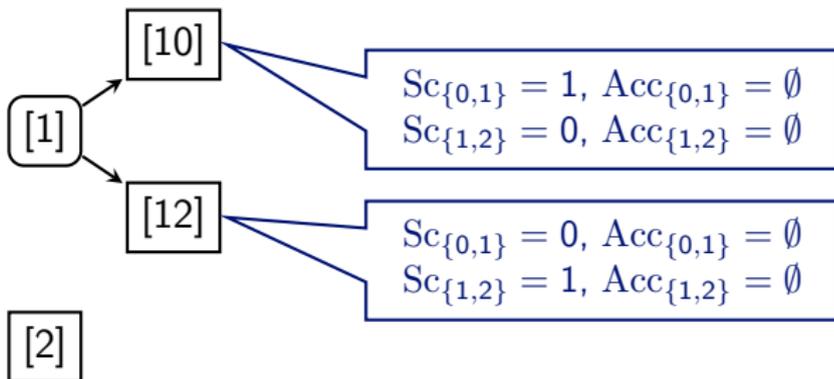
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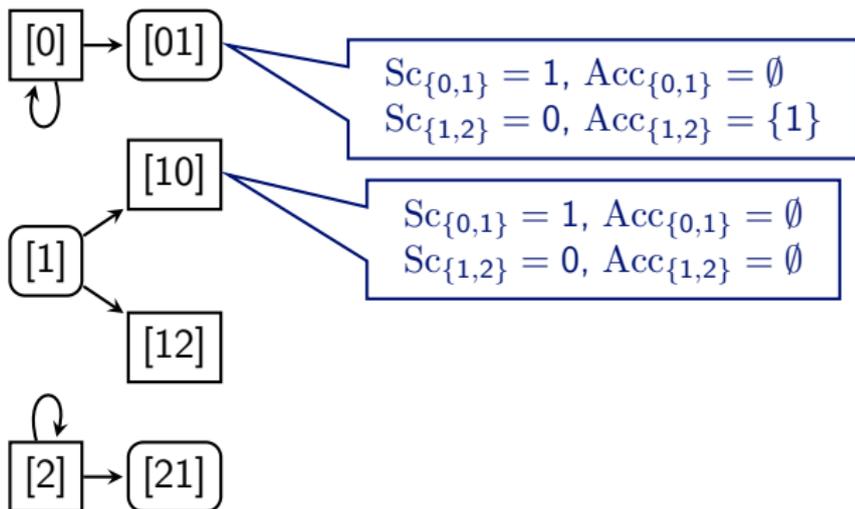


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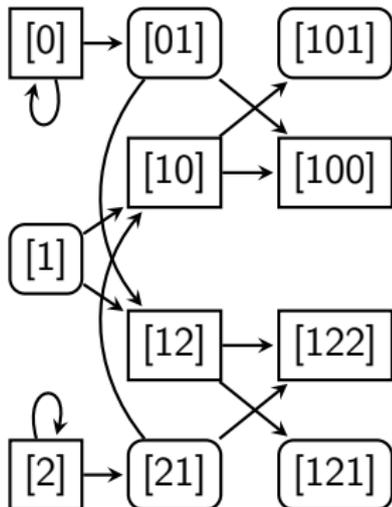
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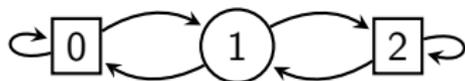
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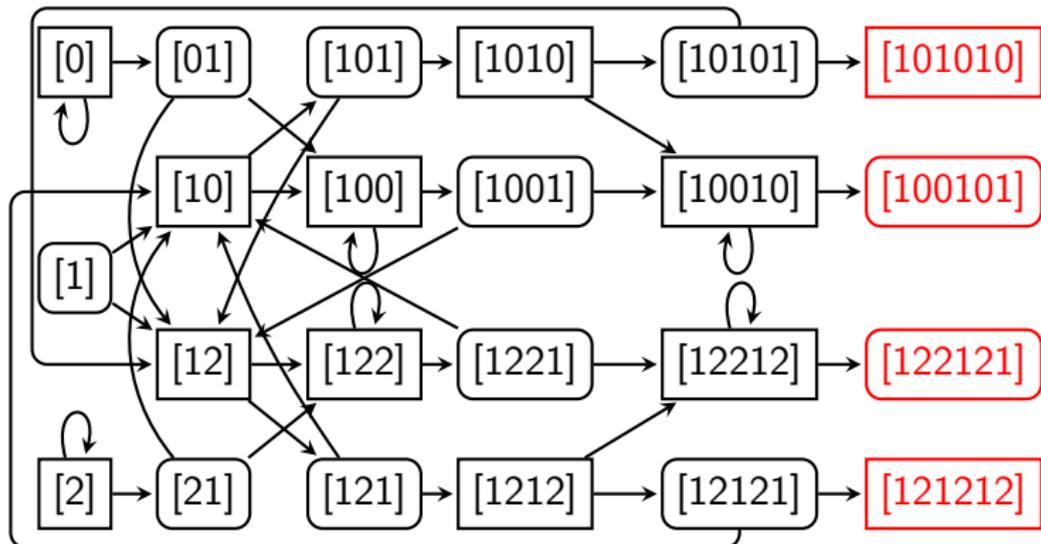


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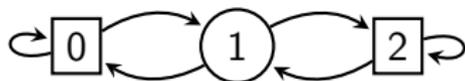


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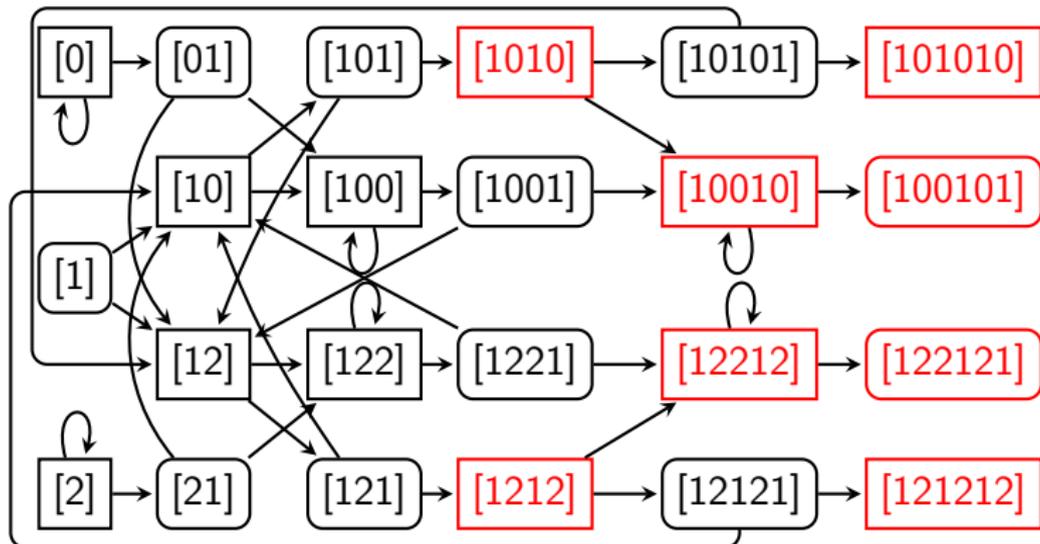
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Theorem (NRZ11)

1. *Player i wins the Muller game from v iff she wins the safety game from $[v]_{=\mathcal{F}_1}$.*
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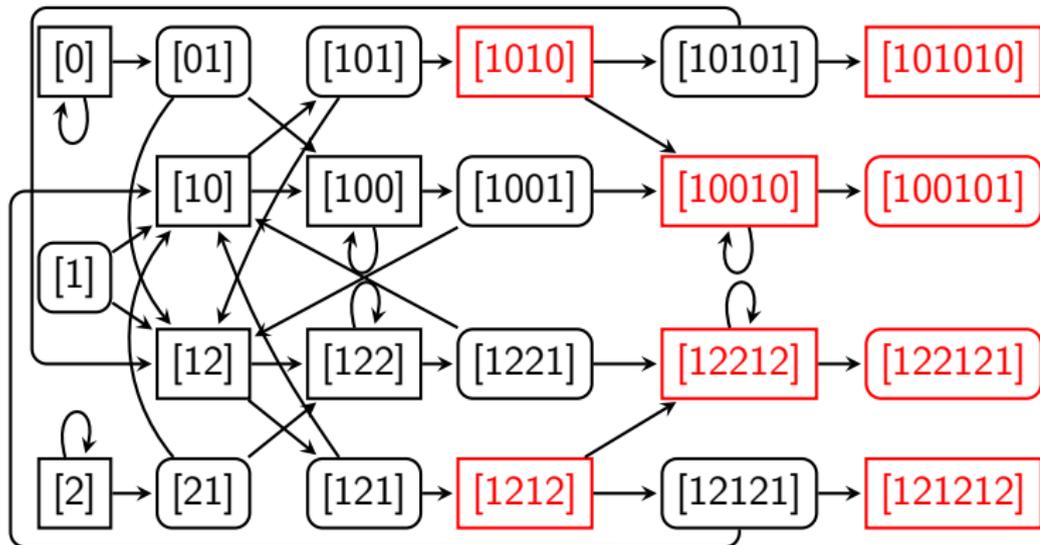
- Size of parity game in LAR-reduction $n!$. But: safety games allow much simpler algorithms.
- 2. does not hold for Player 1.
- Not a reduction in the classical sense: not every play of the Muller game can be mapped to a play in the safety game.

Proof Idea: Safety to Muller



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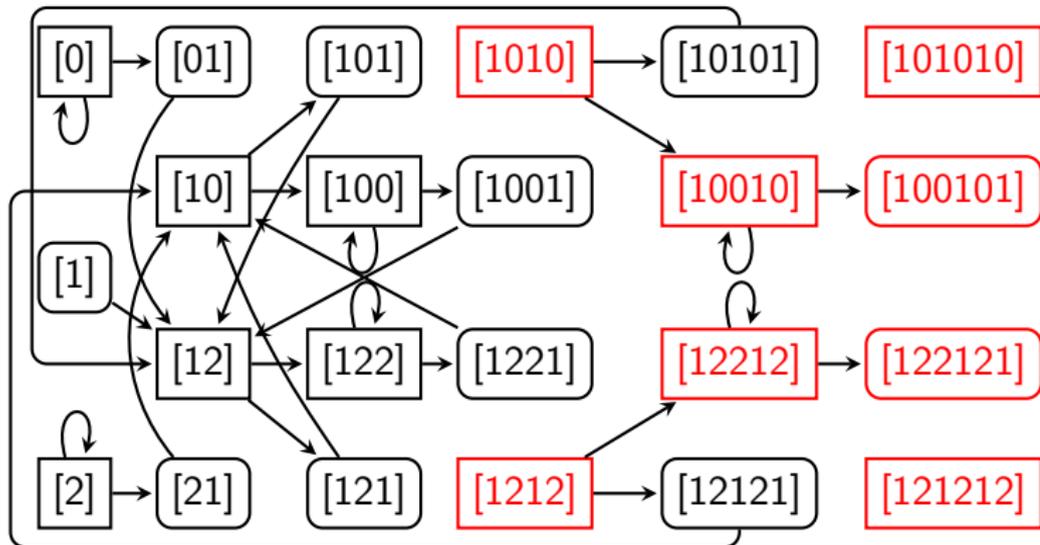


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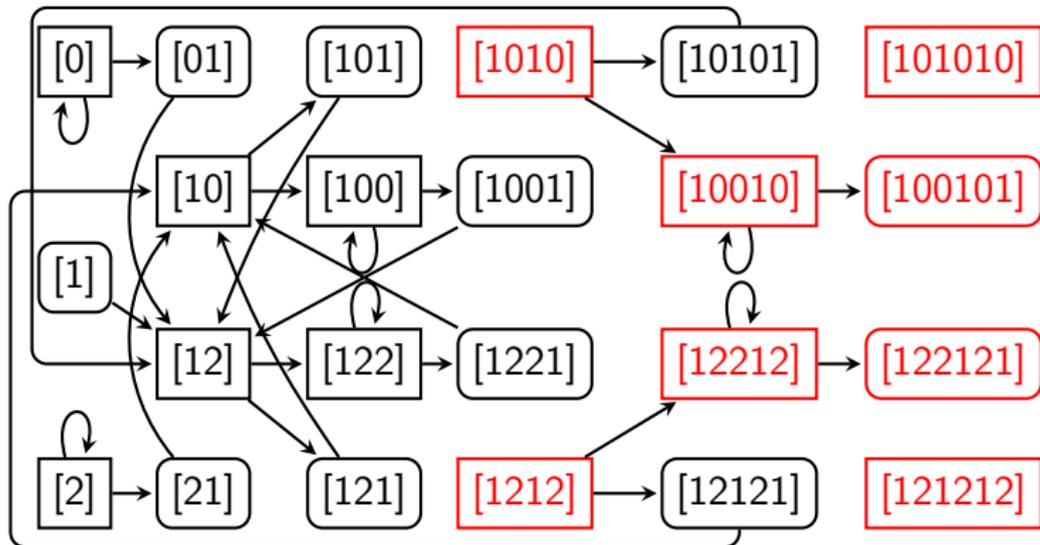
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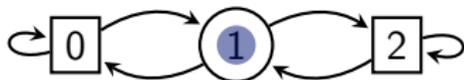
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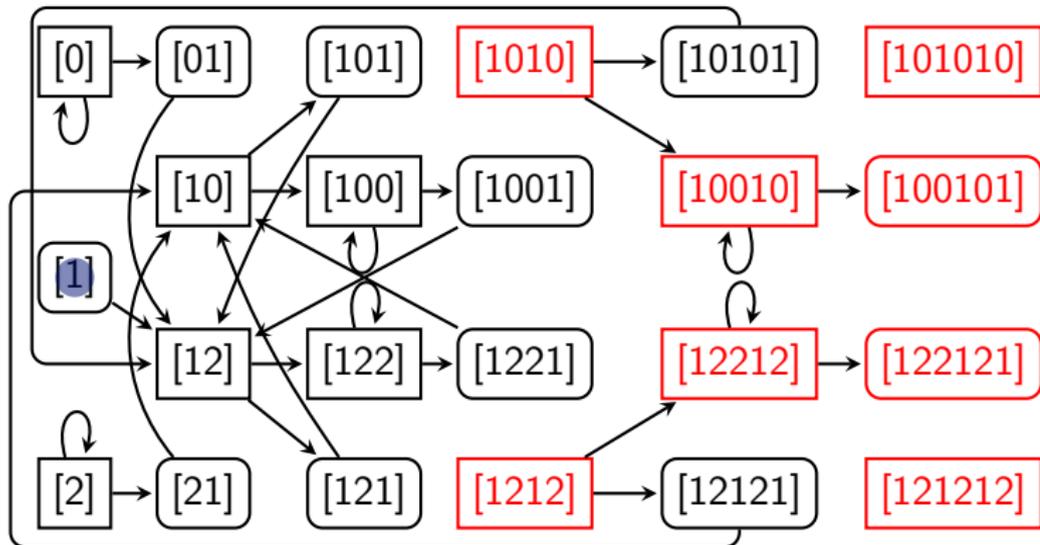
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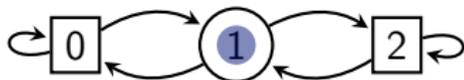
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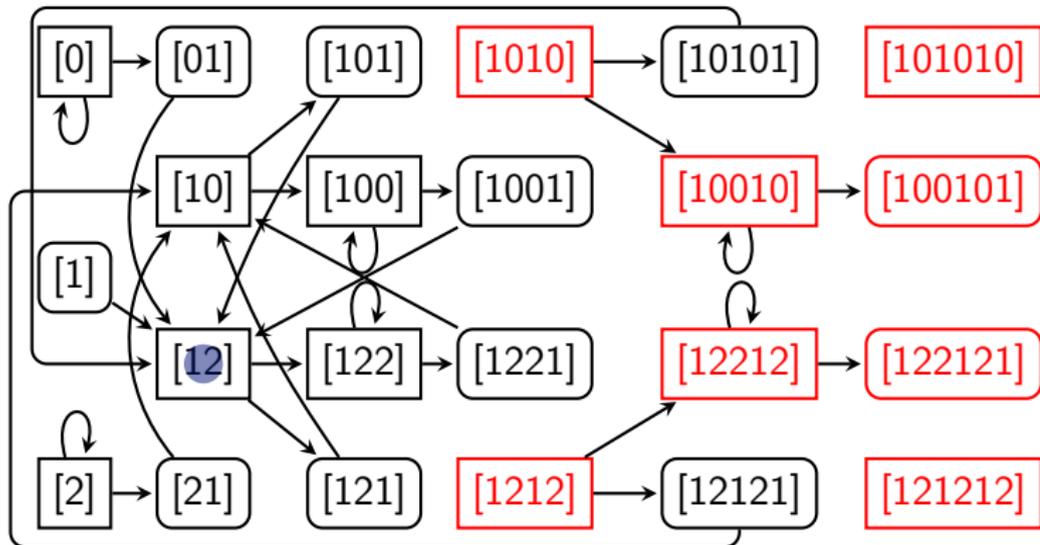
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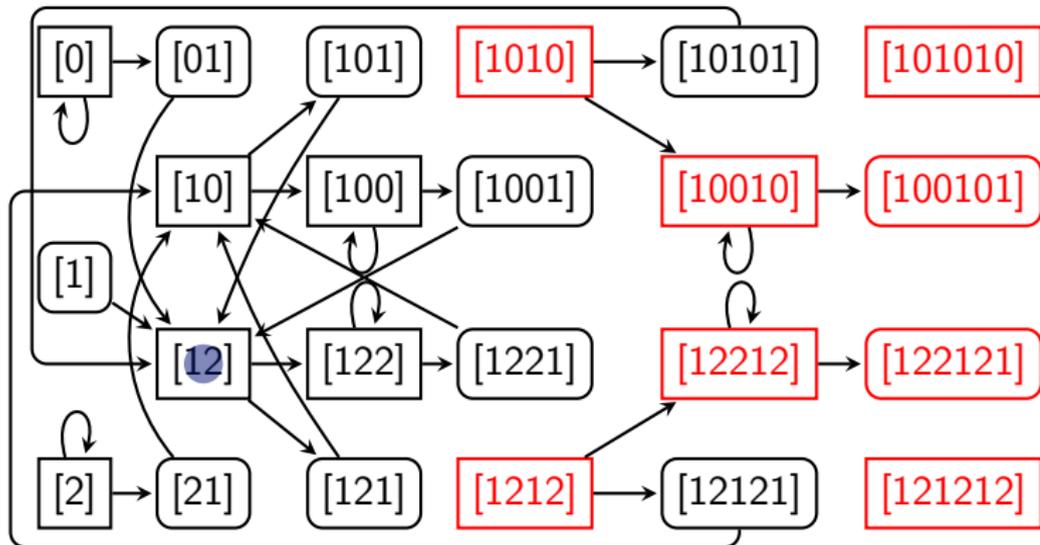
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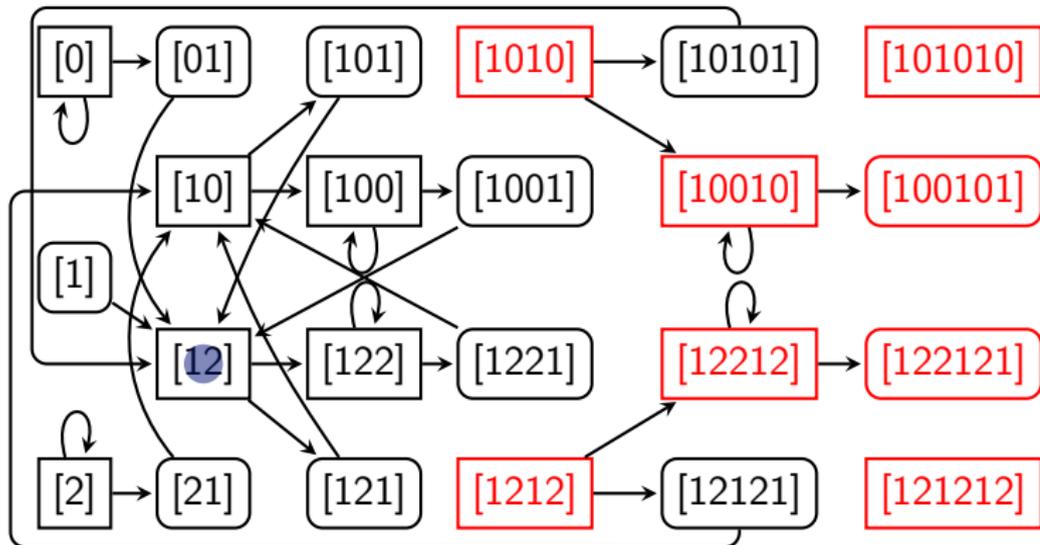
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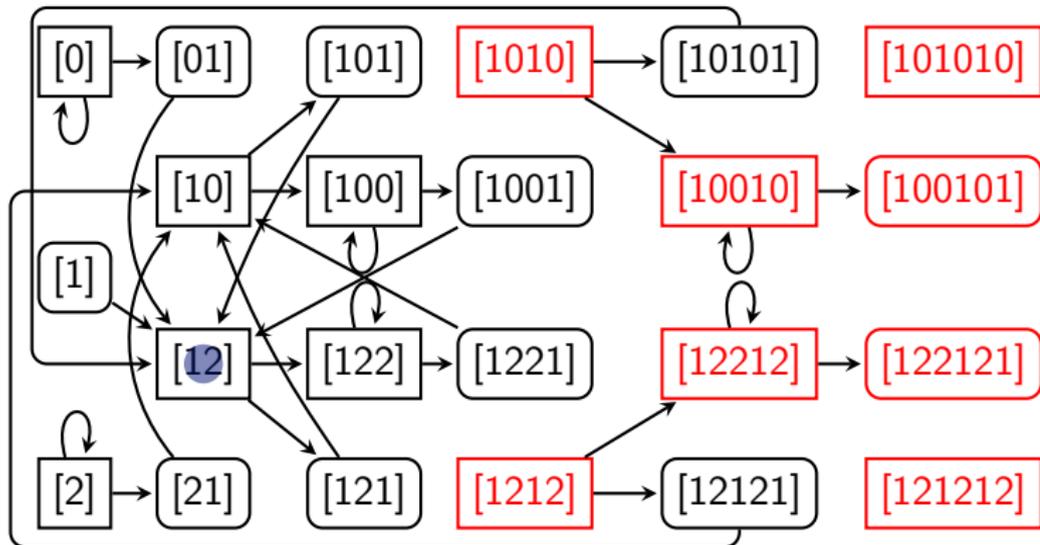
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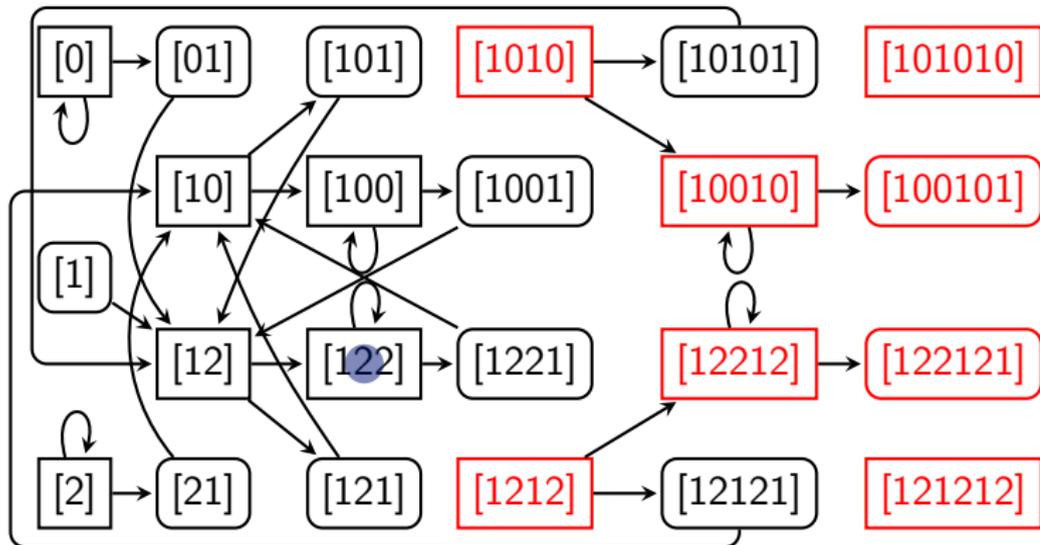
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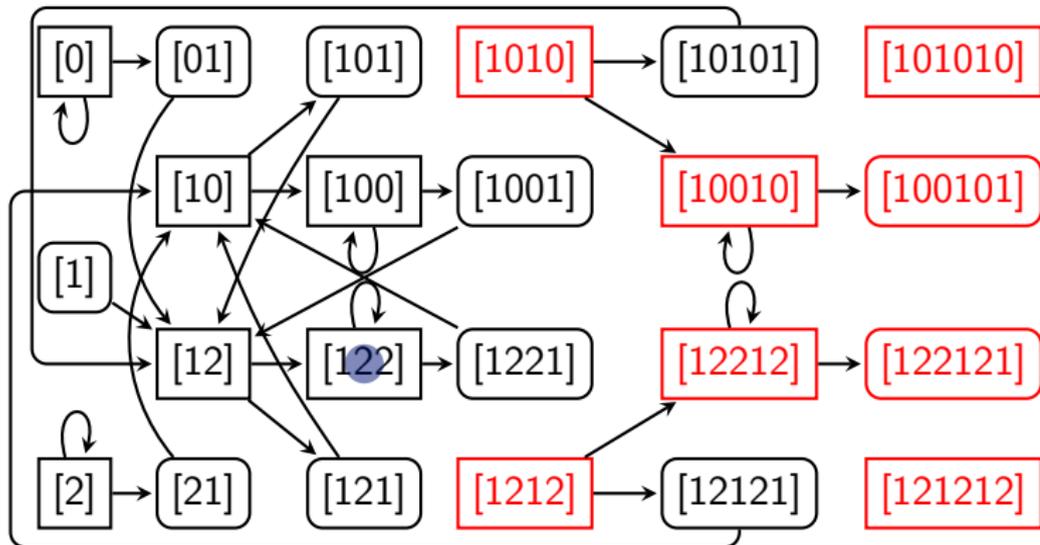
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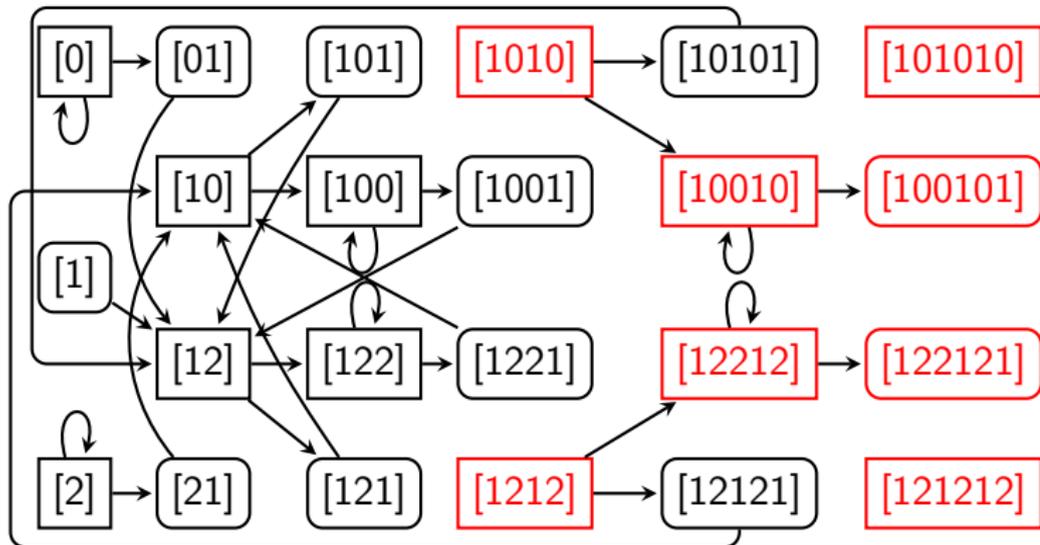
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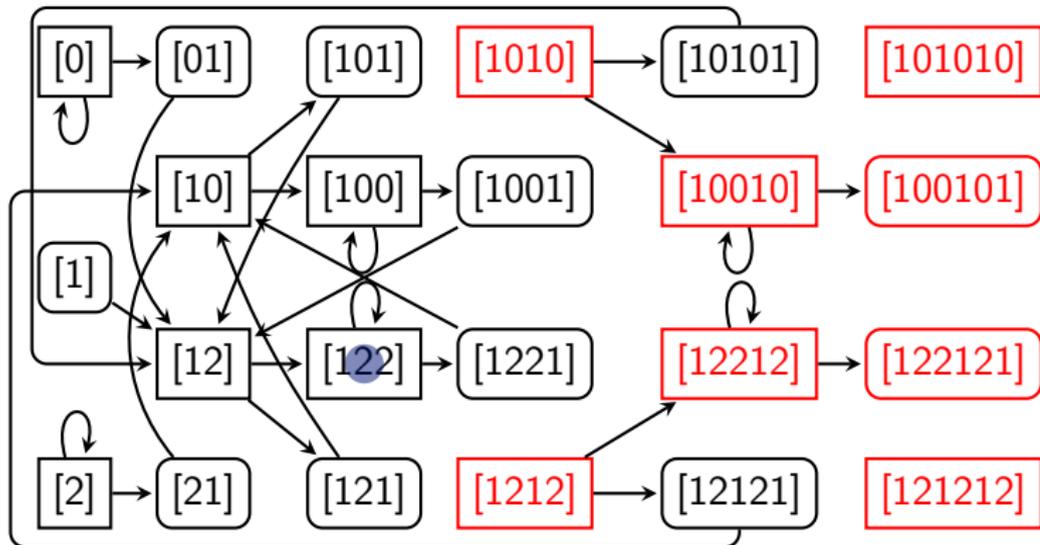
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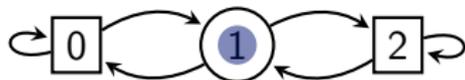
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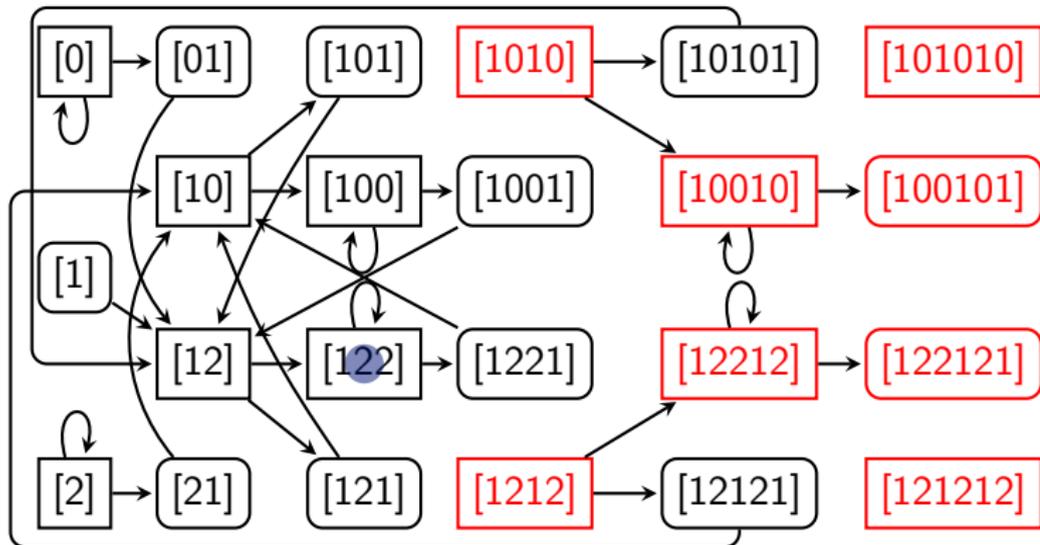
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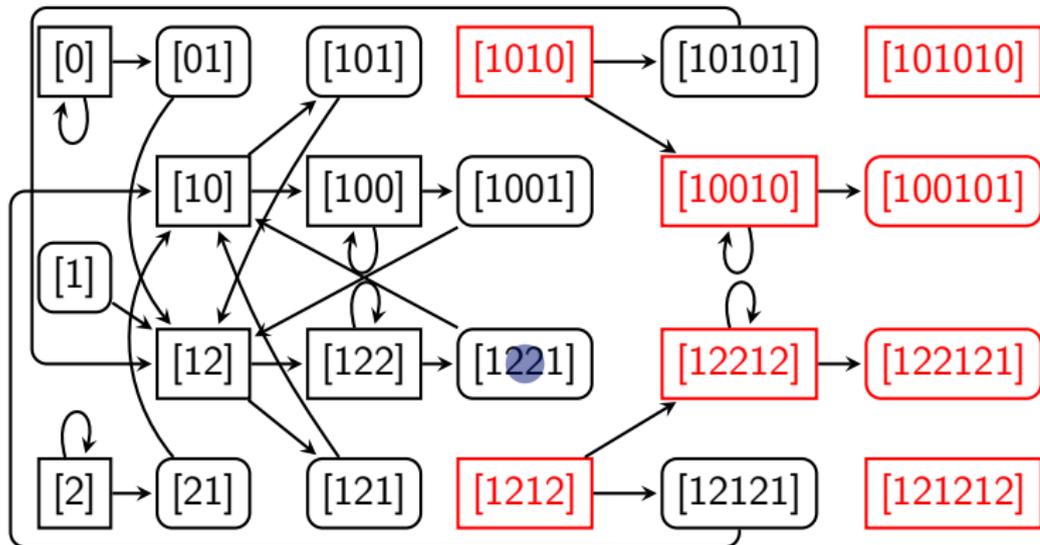
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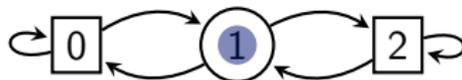
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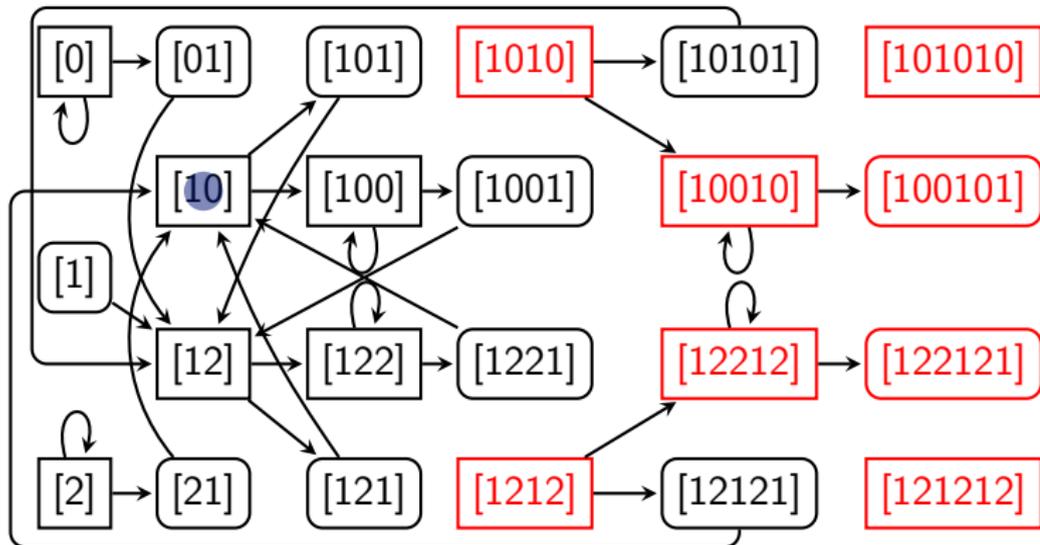
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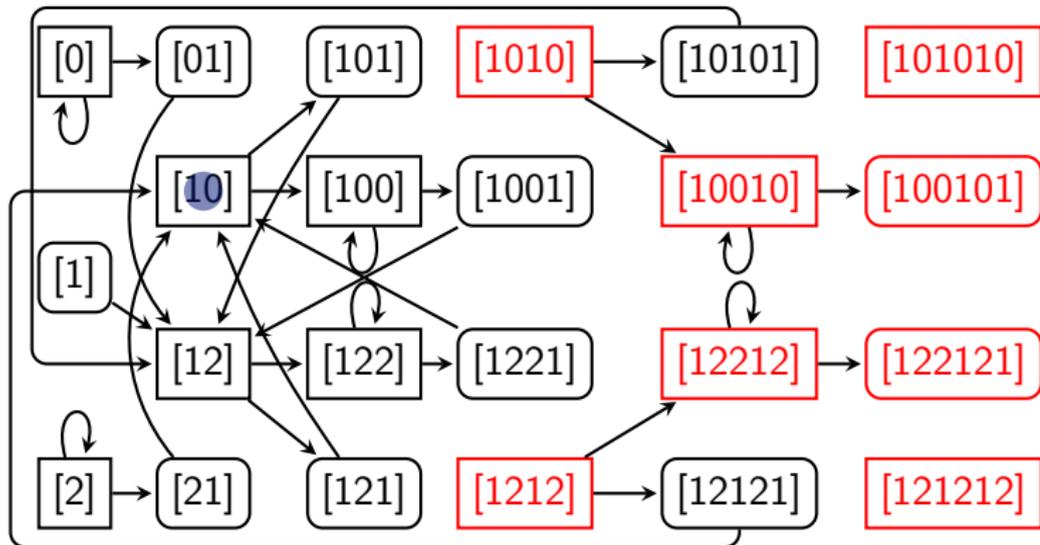
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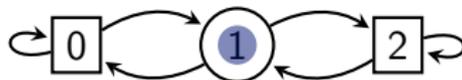
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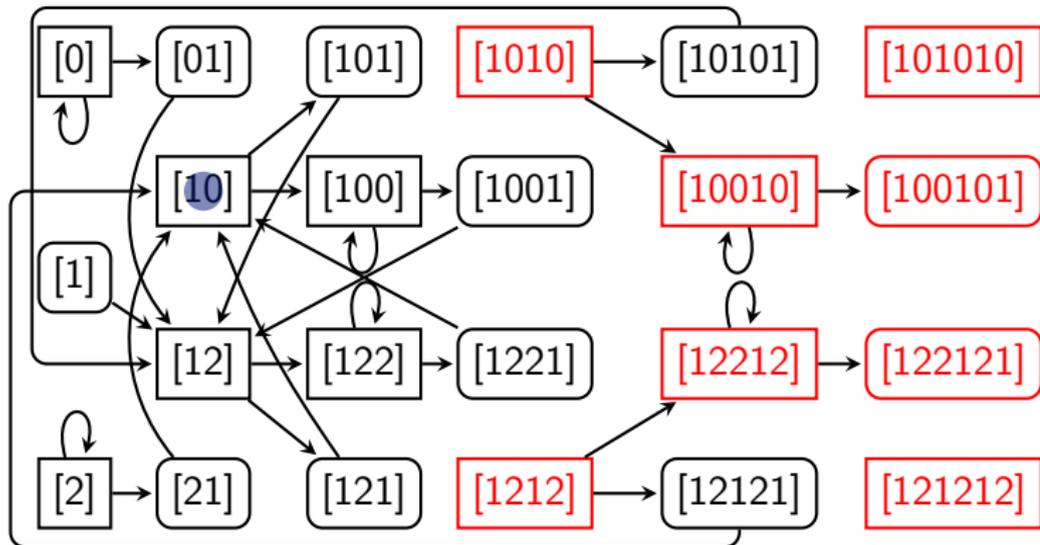
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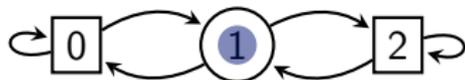
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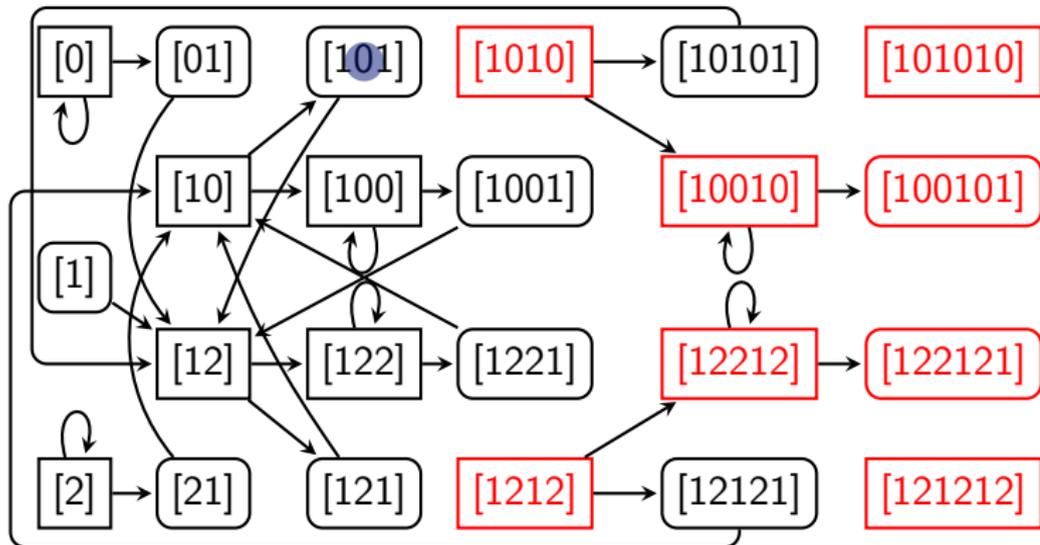


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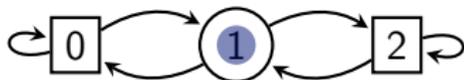


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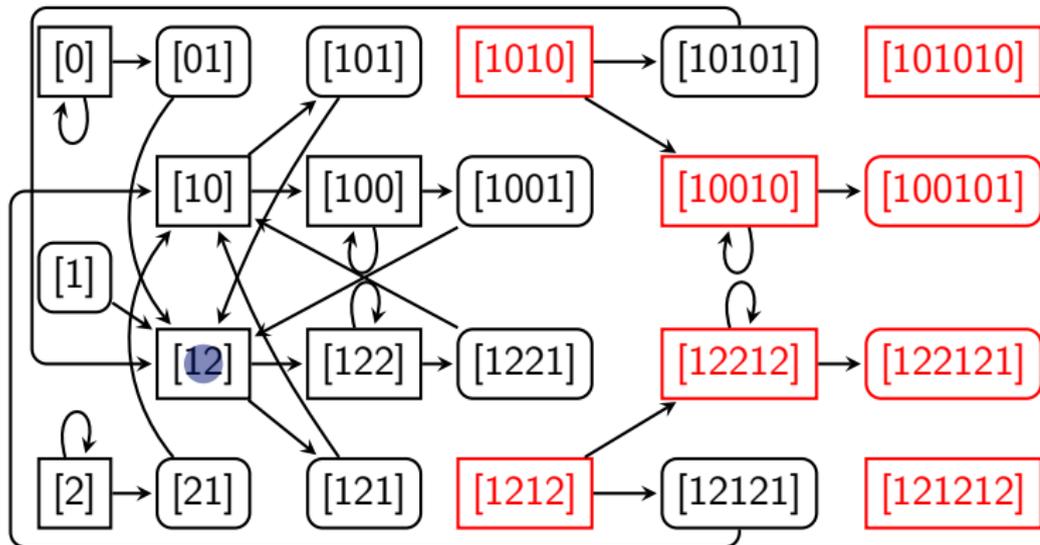
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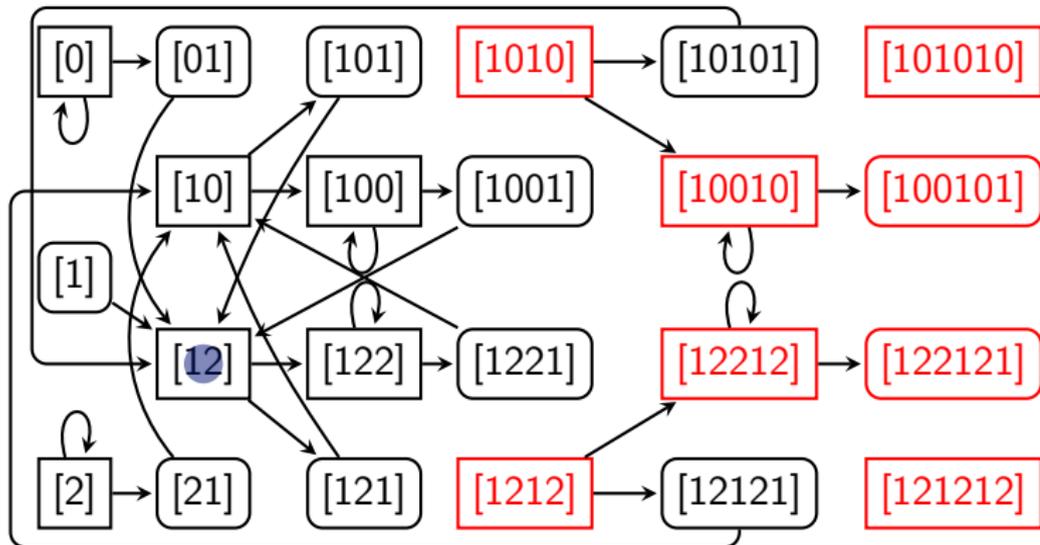
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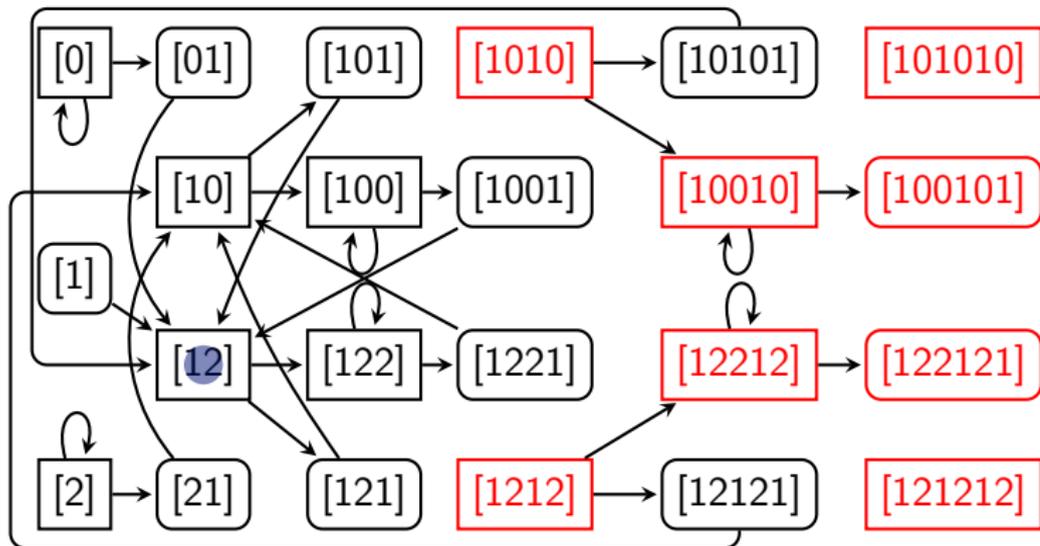
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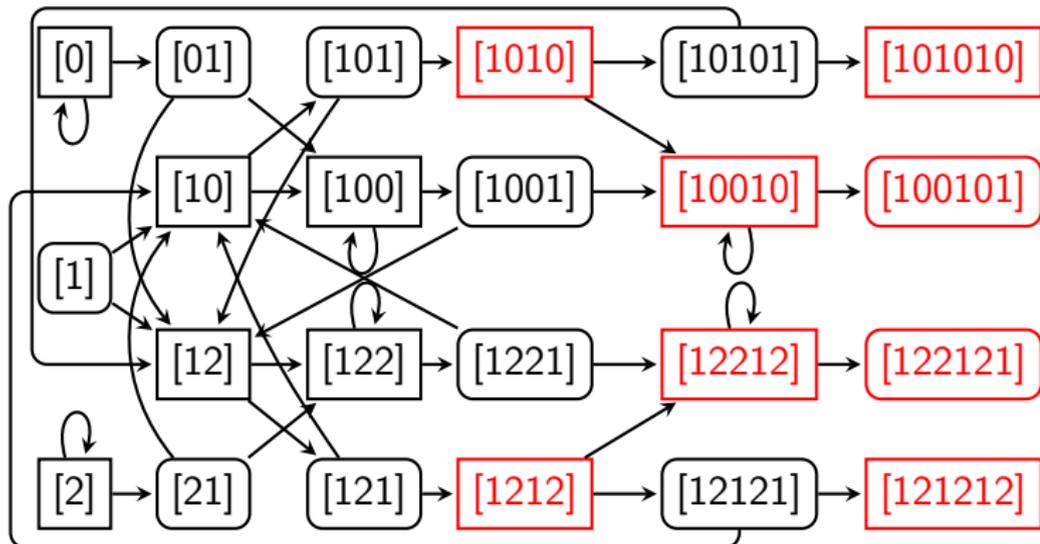
Sc_F for $F \in \mathcal{F}_1$ in Muller game bounded by 2 \Rightarrow winning strategy

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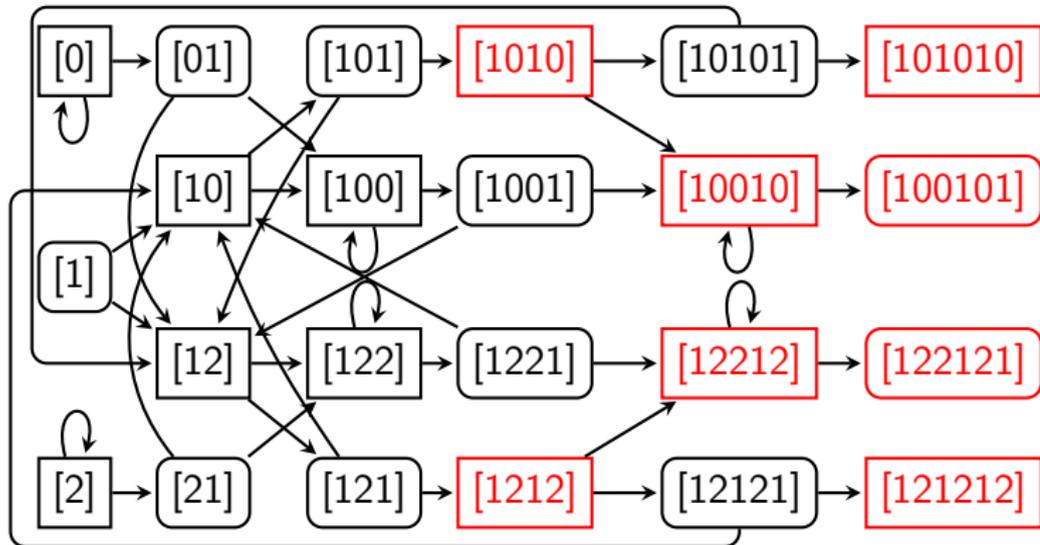


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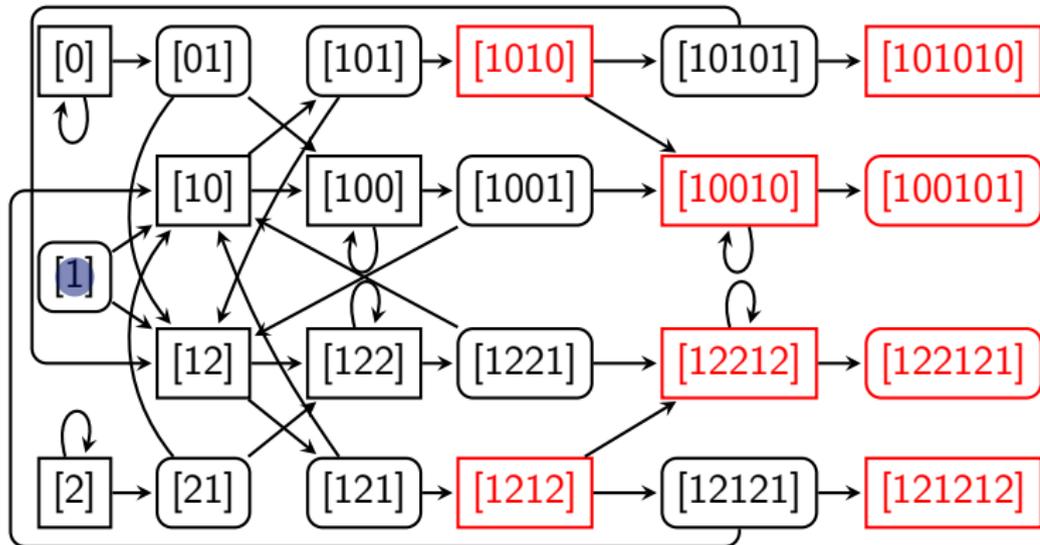
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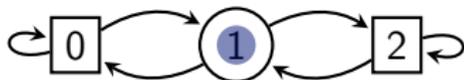
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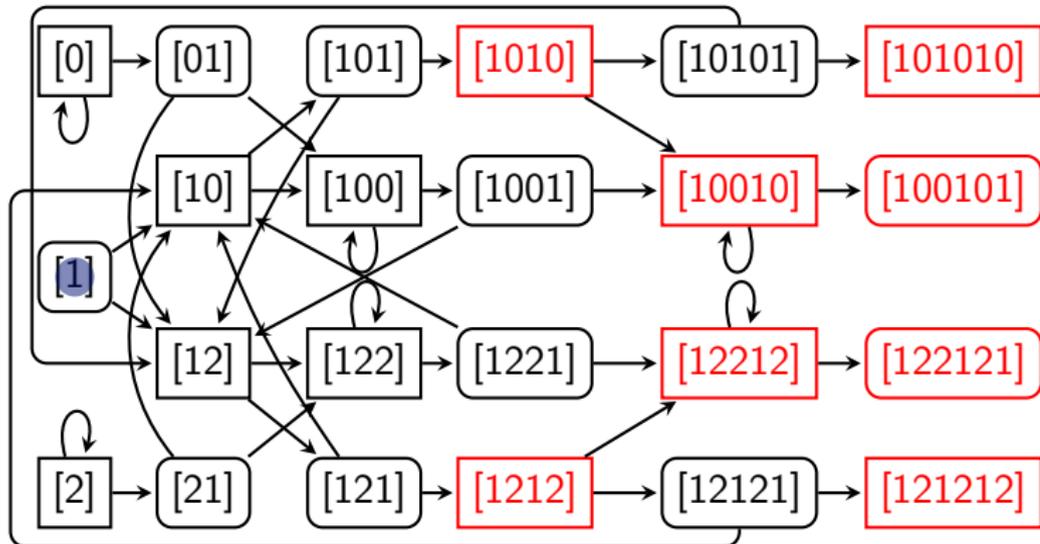
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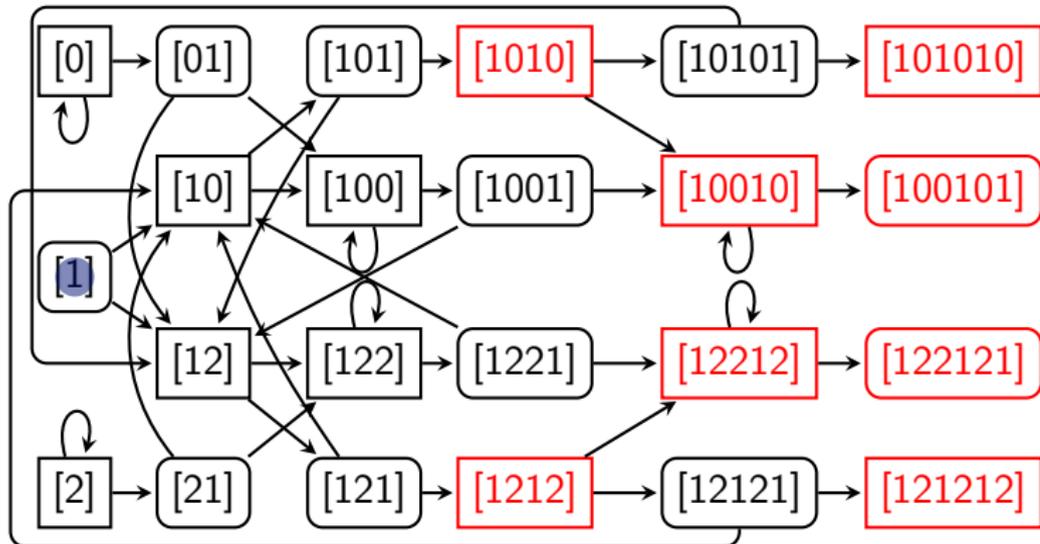


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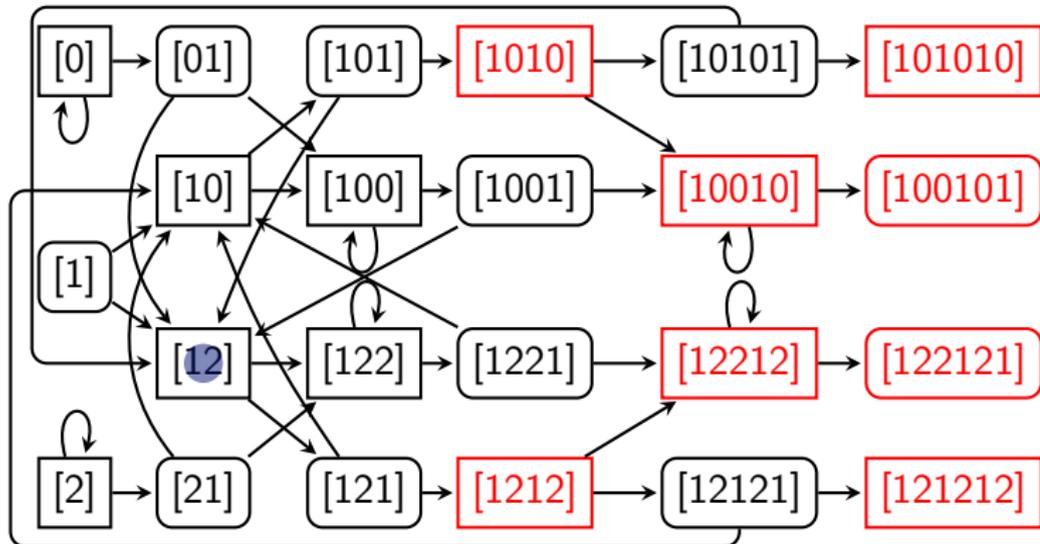
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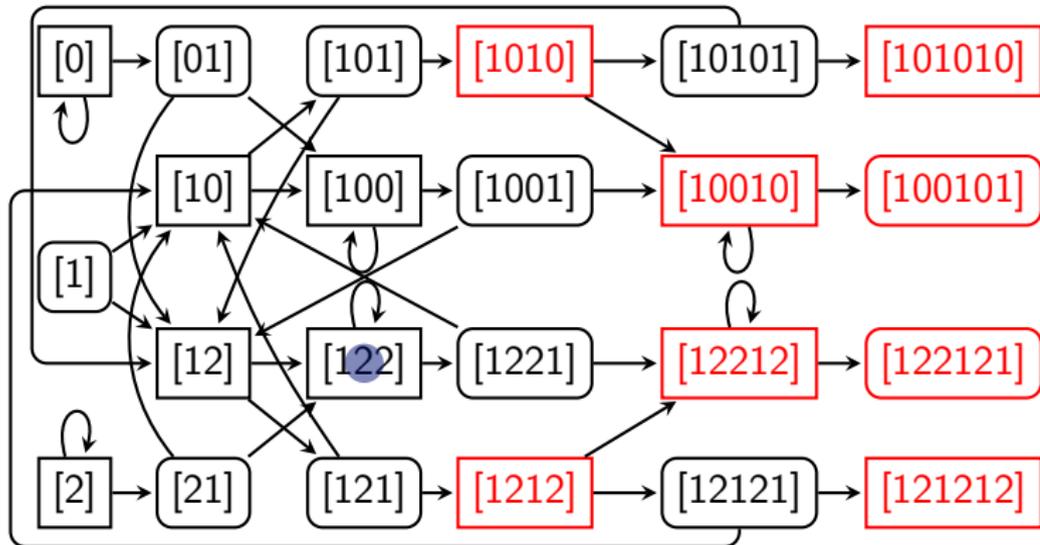
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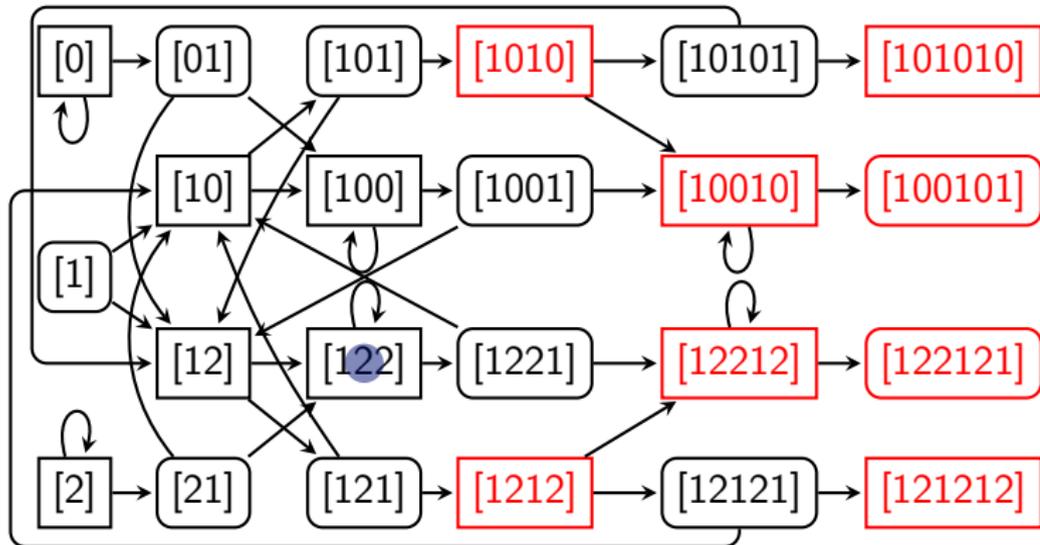
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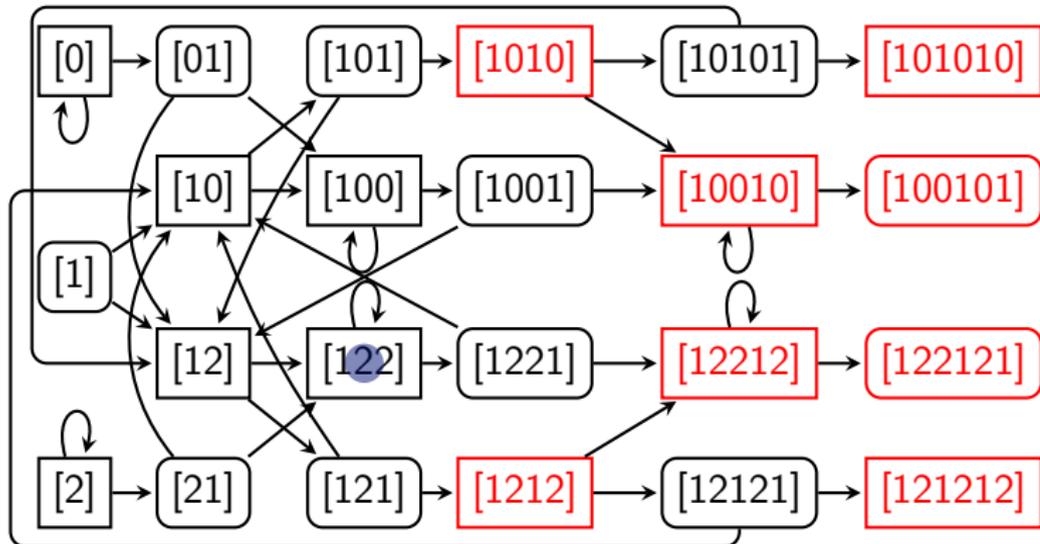
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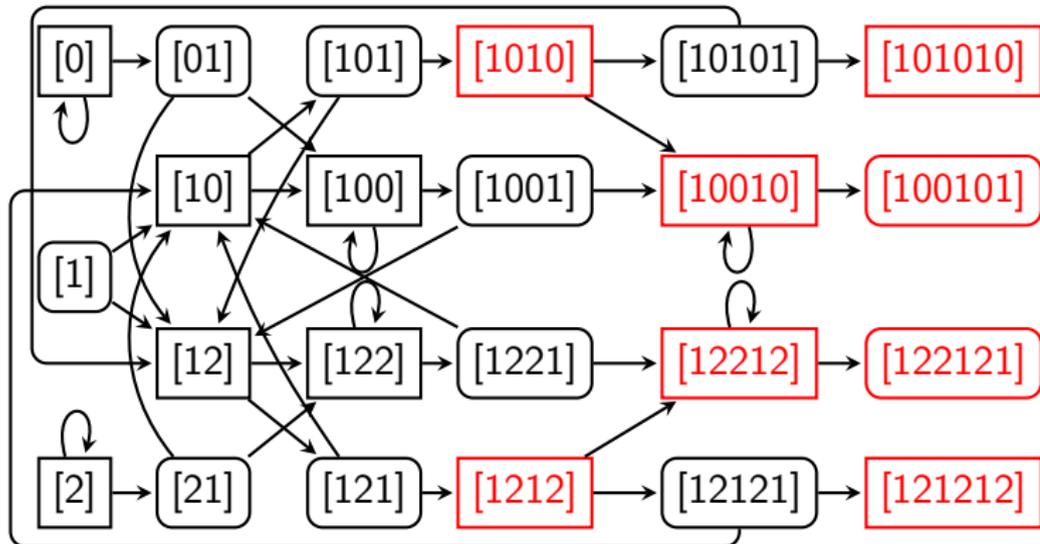
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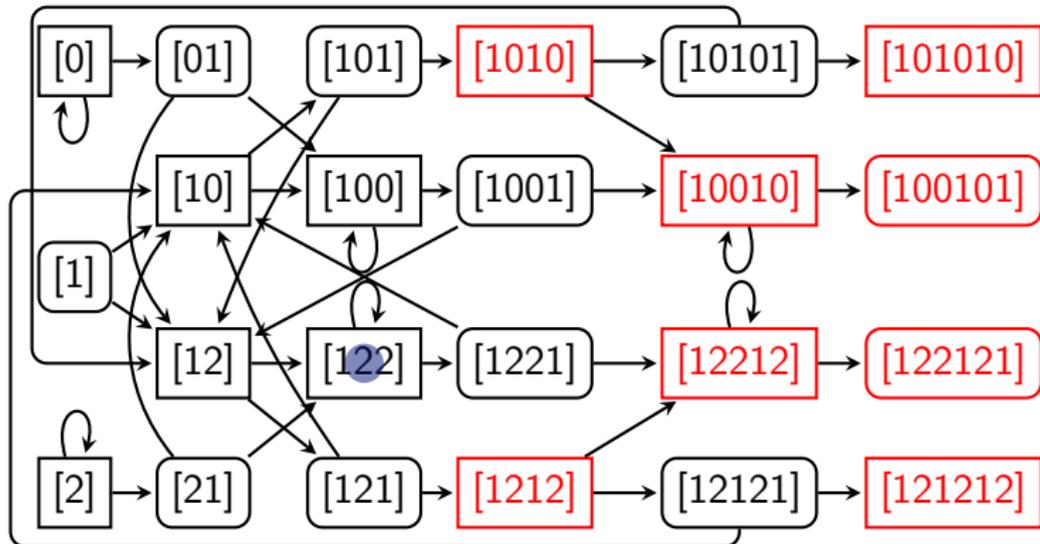
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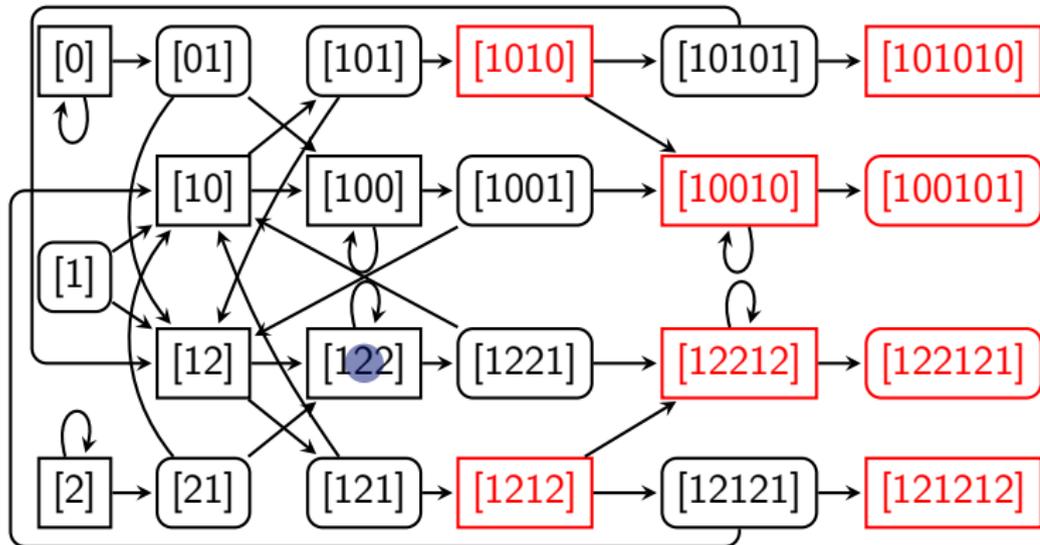
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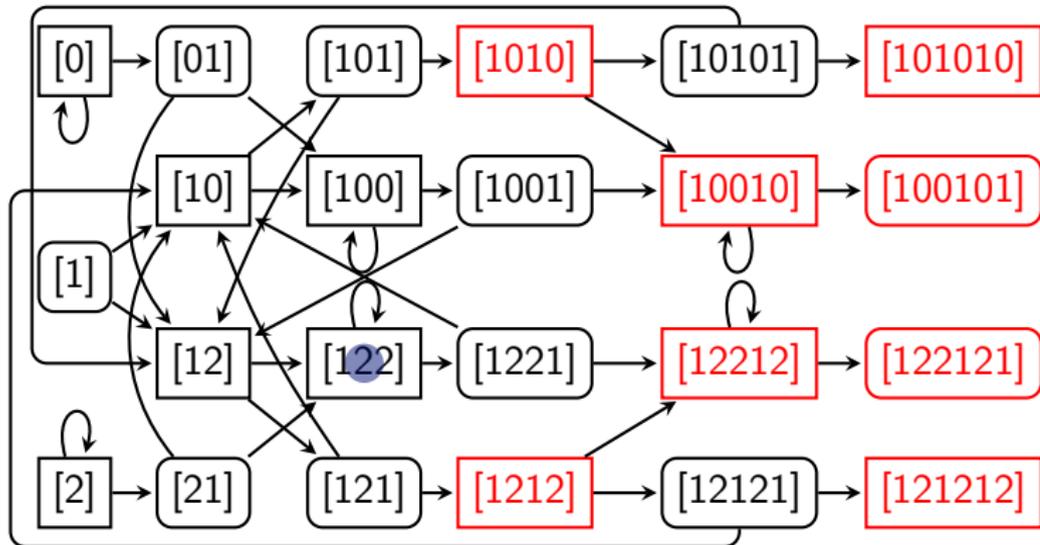
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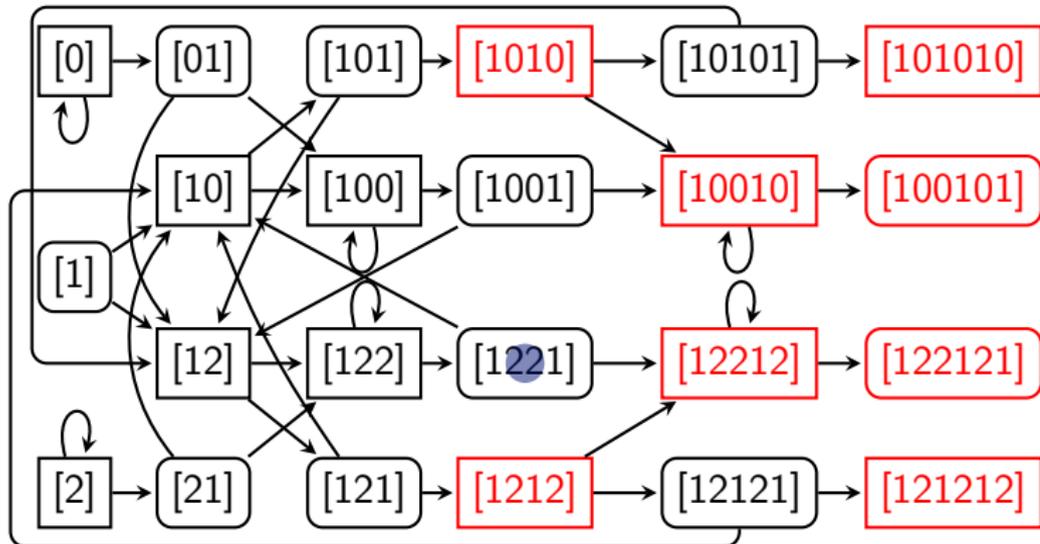
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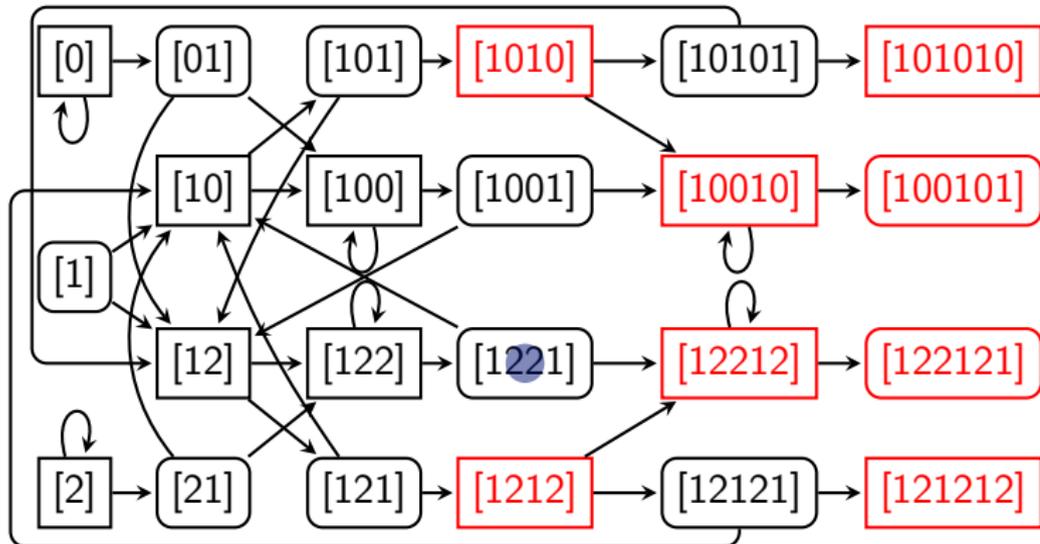


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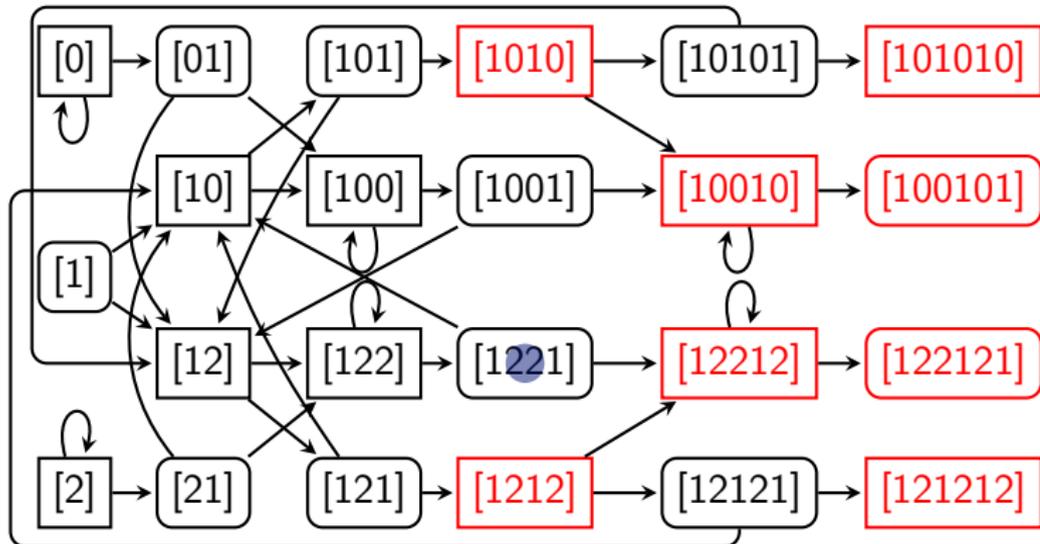


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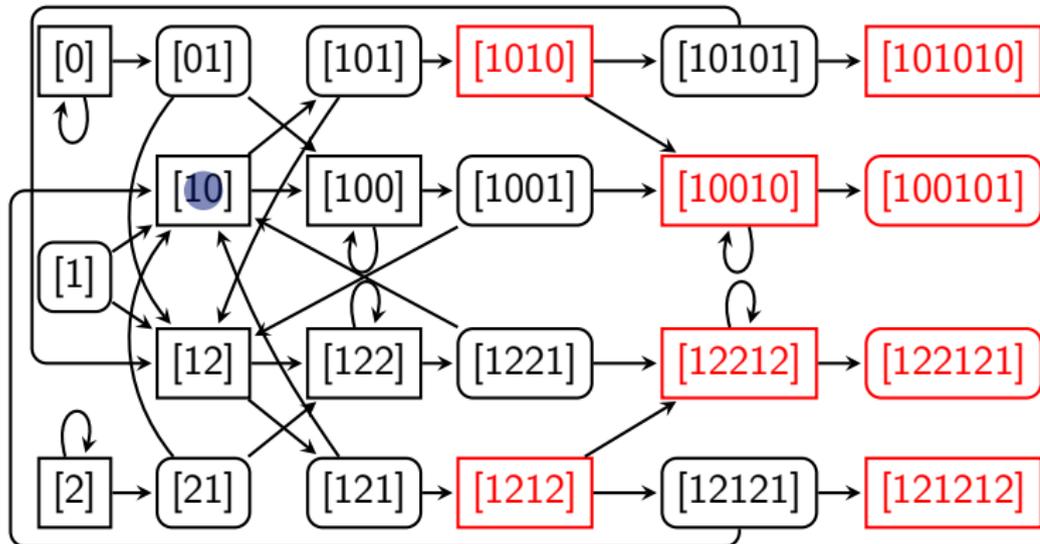


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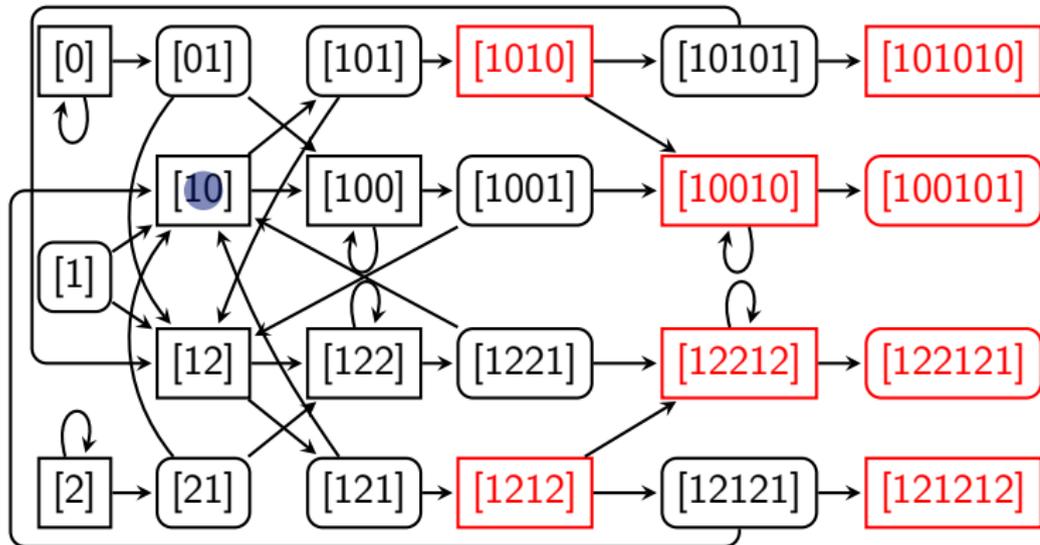
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Score of three is avoidable from every prefix \Rightarrow red vertices never reached \Rightarrow winning strategy.

Outline

1. Playing Muller Games in Finite Time
2. Solving Muller Games by Solving Safety Games
- 3. Conclusion**

Conclusion

You can play Muller games in finite time!

- New algorithm for Muller games: just solve the safety game.
- New memory structure for Muller games: maximal elements of winning region suffice (antichain).
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Ongoing and future work:

- Progress measure algorithm for Muller games?
- Is there a tradeoff between size and quality of a strategy?
- Can you play infinite games in infinite arenas in finite time?