
Down the Borel Hierarchy: Solving Muller Games via Safety Games

Joint work with John Fearnley,
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Introduction

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$$\text{Muller} \xrightarrow{\text{LAR}} \text{parity} \xrightarrow{\text{BJW}} \text{safety}$$

We present a direct construction from Muller to safety with only exponential blowup \Rightarrow new algorithm, new memory structure, and permissive strategies for Muller games.

Muller Games

Muller games $(\mathcal{A}, \mathcal{F}_0, \mathcal{F}_1)$:

- arena \mathcal{A} and partition $(\mathcal{F}_0, \mathcal{F}_1)$ containing the loops of \mathcal{A} .
- Player i wins ρ iff $\text{Inf}(\rho) = \{v \mid \exists^\omega n \text{ s.t. } \rho_n = v\} \in \mathcal{F}_i$.

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Running example



- $\mathcal{F}_0 = \{\{0, 1, 2\}, \{0\}, \{2\}\}$
- $\mathcal{F}_1 = \{\{0, 1\}, \{1, 2\}\}$

Player 0 has a winning strategy from every vertex: alternate between 0 and 2. This requires two memory states.

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Theorem

Muller games are determined with finite-state strategies of size $n!$.

Outline

1. Scoring Functions for Muller Games
2. Solving Muller Games by Solving Safety Games
3. Conclusion

Scoring Functions

Let $F \subseteq V$, $F \neq \emptyset$.

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$$\text{Sc}_F(v) = \begin{cases} 1 & \text{if } F = \{v\}, \\ 0 & \text{otherwise,} \end{cases}$$

and

$$\text{Acc}_F(v) = \begin{cases} \emptyset & \text{if } F = \{v\}, \\ F \cap \{v\} & \text{otherwise.} \end{cases}$$

Scoring Functions

Let $F \subseteq V$, $F \neq \emptyset$. For $v \in V$ and $w \in V^+$ define

$$\text{Sc}_F(wv) = \begin{cases} 0 & \text{if } v \notin F, \\ \text{Sc}_F(w) & \text{if } v \in F \wedge \text{Acc}_F(w) \neq (F \setminus \{v\}), \\ \text{Sc}_F(w) + 1 & \text{if } v \in F \wedge \text{Acc}_F(w) = (F \setminus \{v\}), \end{cases}$$

and

$$\text{Acc}_F(wv) = \begin{cases} \emptyset & \text{if } v \notin F, \\ \text{Acc}_F(w) \cup \{v\} & \text{if } v \in F \wedge \text{Acc}_F(w) \neq (F \setminus \{v\}), \\ \emptyset & \text{if } v \in F \wedge \text{Acc}_F(w) = (F \setminus \{v\}). \end{cases}$$

Scoring Functions cont'd

- $Sc_F(w)$: maximal $k \in \mathbb{N}$ such that F is visited k times since last vertex in $V \setminus F$ (reset).
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Example:

w		0	0	1	1	0	0	1	2
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$Sc_{\{0,1,2\}}$	0	0	0	0	0	0	0	1
$Acc_{\{0,1,2\}}$	{0}	{0}	{0, 1}	{0, 1}	{0, 1}	{0, 1}	{0, 1}	\emptyset

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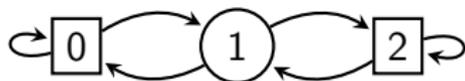
Example:

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$Acc_{\{0,1,2\}}$	{0}	{0}	{0, 1}	{0, 1}	{0, 1}	{0, 1}	{0, 1}	\emptyset

Remark

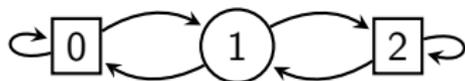
$$F = \text{Inf}(\rho) \Leftrightarrow \liminf_{n \rightarrow \infty} Sc_F(\rho_0 \cdots \rho_n) = \infty$$

Two Examples



- $\mathcal{F}_0 = \{\{0, 1, 2\}, \{0\}, \{2\}\}$
- $\mathcal{F}_1 = \{\{0, 1\}, \{1, 2\}\}$

Two Examples



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Losing player (Player 1) can enforce score of two:

Two Examples



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Two Examples



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Losing player (Player 1) can enforce score of two:

$$1 \rightarrow 2 \quad (\text{w.l.o.g.})$$

Two Examples



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Losing player (Player 1) can enforce score of two:

$$1 \rightarrow 2 \rightarrow 2$$

Two Examples



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Losing player (Player 1) can enforce score of two:

$$1 \rightarrow 2 \rightarrow 2 \rightarrow 1$$

Two Examples



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Losing player (Player 1) can enforce score of two:

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Two Examples

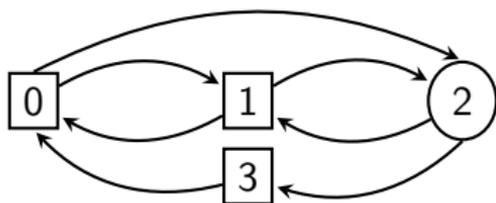


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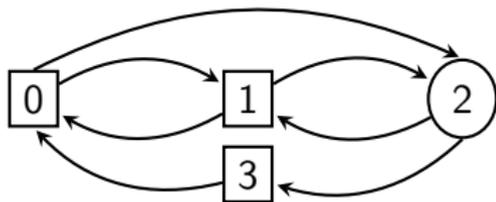


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Two Examples

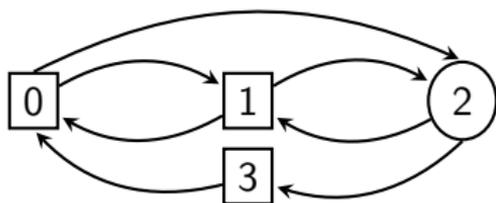


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Losing player (Player 1) is the first to reach a score of two:

3

Two Examples

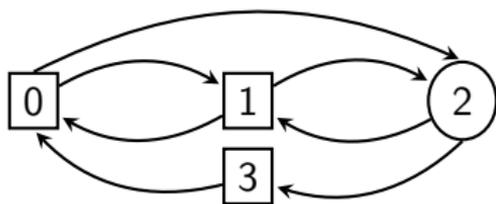


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Losing player (Player 1) is the first to reach a score of two:

$$3 \rightarrow 0$$

Two Examples

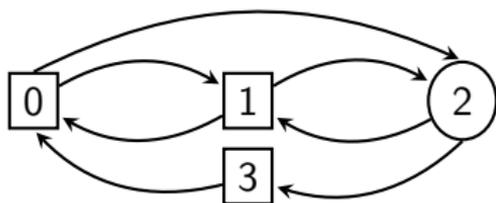


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Two Examples

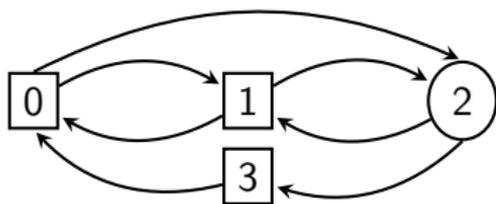


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Losing player (Player 1) is the first to reach a score of two:

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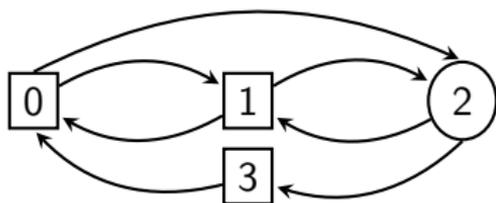


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Two Examples

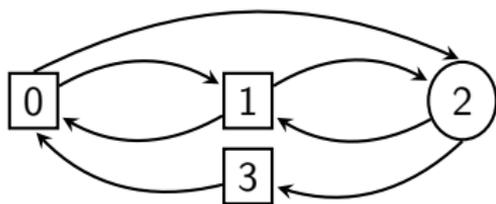


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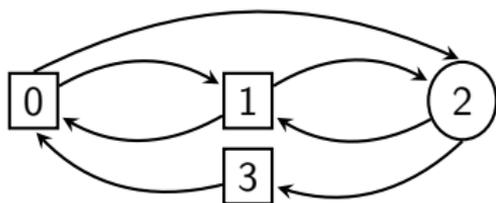


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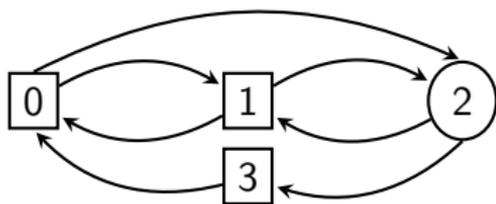


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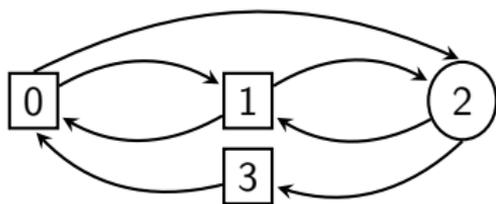


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Two Examples

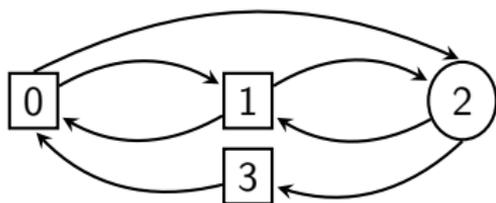


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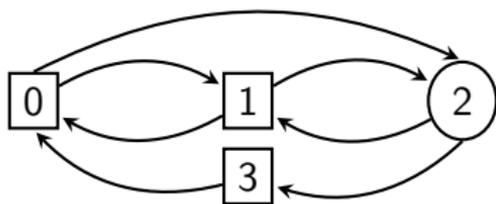


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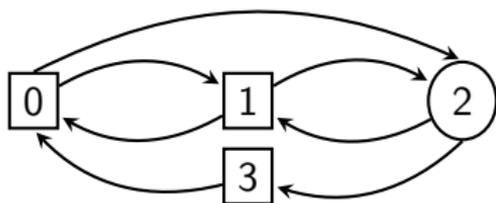


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Losing player (Player 1) is the first to reach a score of two:

$$3 \rightarrow 0 \rightarrow 2 \begin{cases} \rightarrow 1 \rightarrow 0 \rightarrow 1 \rightarrow 2 \\ \rightarrow 3 \rightarrow 0 \rightarrow 2 \end{cases} \quad \leftarrow \boxed{\text{Sc}_{\{0,1,2\}} = 2}$$
$$3 \rightarrow 0 \rightarrow 2 \quad \leftarrow \boxed{\text{Sc}_{\{0,2,3\}} = 2}$$

Lemma (FZ10)

On her winning region, Player i can prevent her opponent from ever reaching a score of 3 for every set $F \in \mathcal{F}_{1-i}$.

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Corollary

Two “reductions”: Muller game to..

- 1. ..reachability game on unraveling up to score 3: yields winning regions, but no winning strategies.*
- 2. ..safety game: see next slides: yields winning regions and one winning strategy.*

Remember: winning regions and one winning strategy is the best we can hope for.

Outline

1. Scoring Functions for Muller Games
- 2. Solving Muller Games by Solving Safety Games**
3. Conclusion

“Reducing“ Muller Games to Safety Games



- $\mathcal{F}_0 = \{\{0, 1, 2\}, \{0\}, \{2\}\}$
- $\mathcal{F}_1 = \{\{0, 1\}, \{1, 2\}\}$

“Reducing“ Muller Games to Safety Games



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Idea: track of Player 1's scores and avoid $\text{Sc}_F = 3$ for $F \in \mathcal{F}_1$.

“Reducing“ Muller Games to Safety Games



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Idea: track of Player 1's scores and avoid $\text{Sc}_F = 3$ for $F \in \mathcal{F}_1$.

- Ignore scores of Player 0.
- Identify plays having the same scores and accumulators for Player 1: $w =_{\mathcal{F}_1} w'$ iff $\text{last}(w) = \text{last}(w')$ and for all $F \in \mathcal{F}_1$:

$$\text{Sc}_F(w) = \text{Sc}_F(w') \text{ and } \text{Acc}_F(w) = \text{Acc}(w')$$

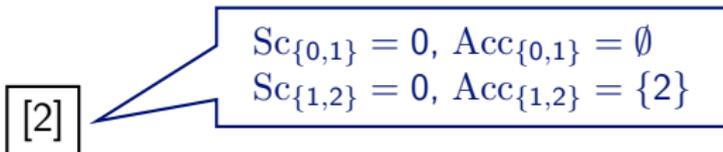
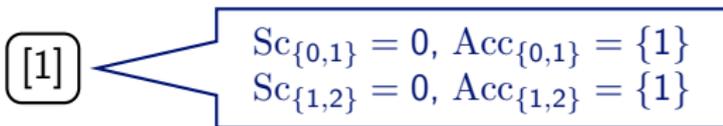
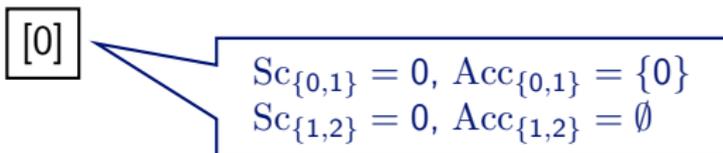
- Build $=_{\mathcal{F}_1}$ -quotient of unravelling up to score 3 for Player 1.
- Winning condition for Player 0: avoid $\text{Sc}_F = 3$ for all $F \in \mathcal{F}_1$.

“Reducing“ Muller Games to Safety Games



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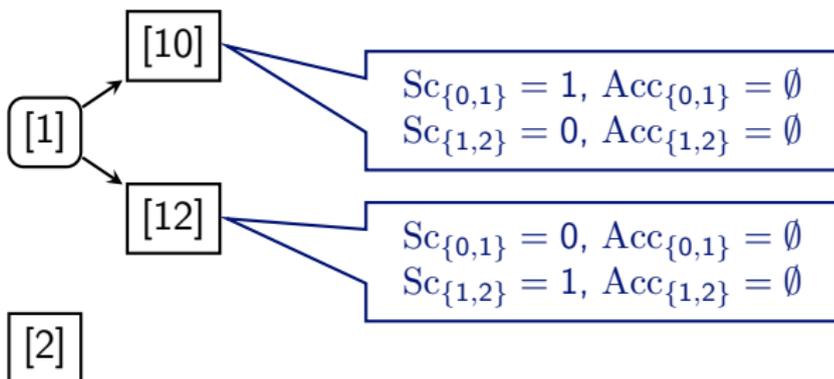
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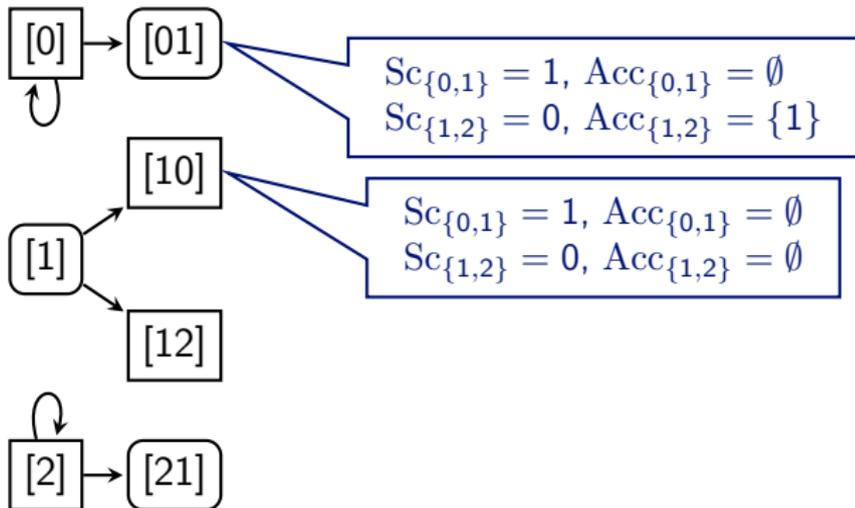


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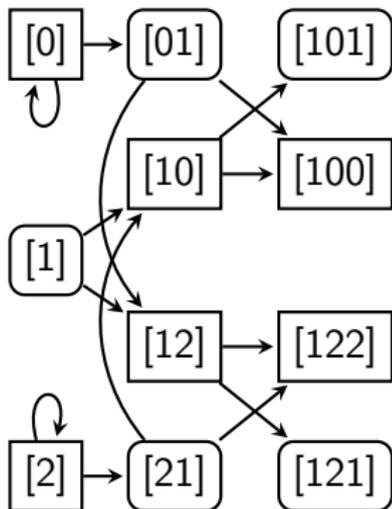
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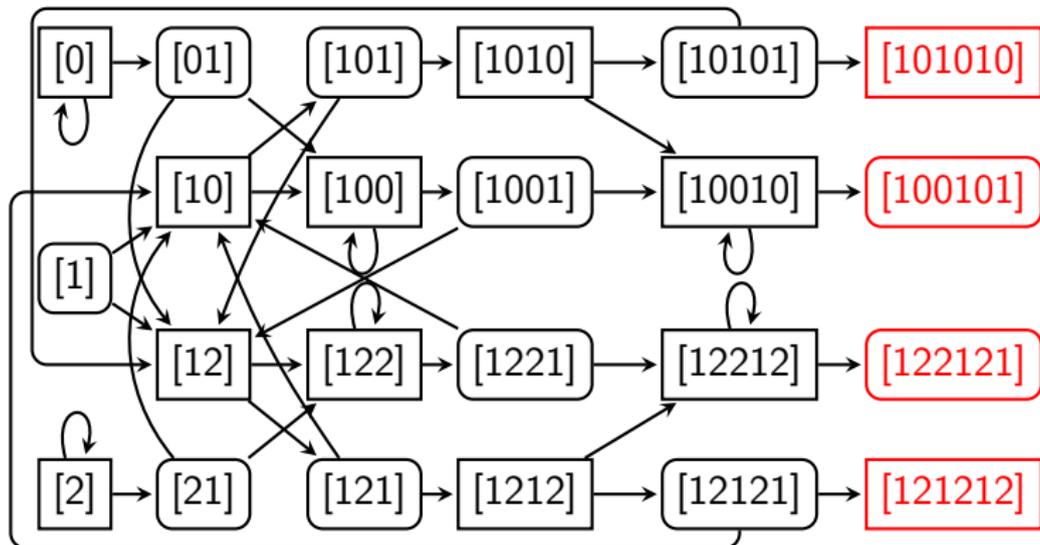


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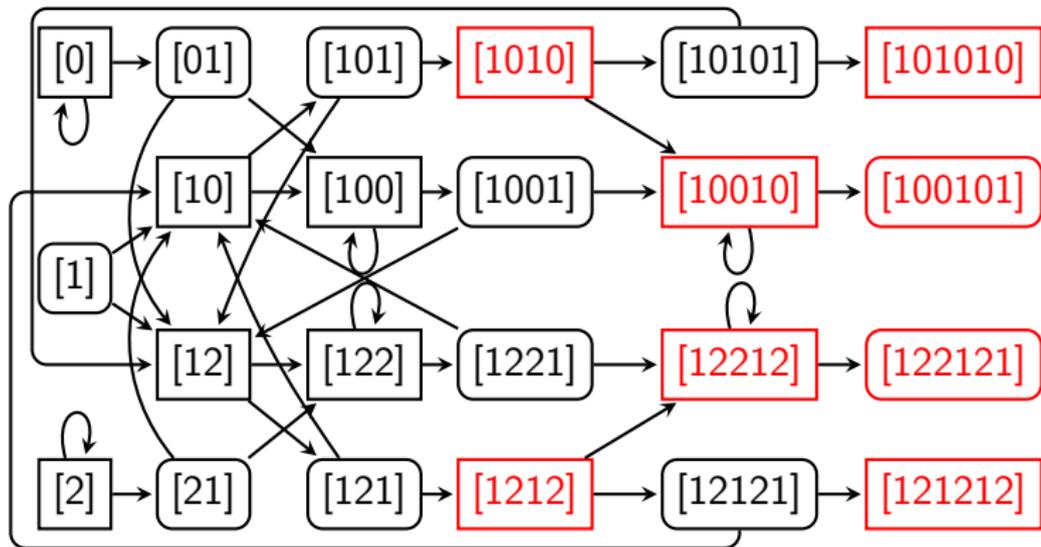


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Theorem (NRZ11)

1. *Player i wins the Muller game from v iff she wins the safety game from $[v]_{=\mathcal{F}_1}$.*
2. *Player 0's winning region in the safety game can be turned into finite-state winning strategy for her in the Muller game.*
3. *Size of the safety game $(n!)^3$.*

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Remarks:

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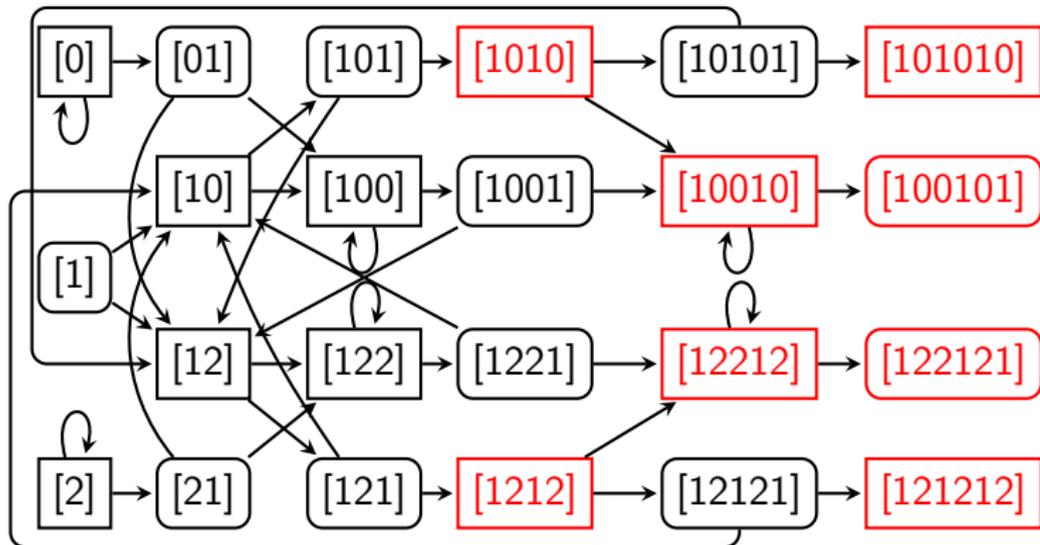
- Size of parity game in LAR-reduction $n!$. But: safety games allow much simpler algorithms.
- 2. does not hold for Player 1.
- Not a reduction in the classical sense: not every play of the Muller game can be mapped to a play in the safety game.

The Proof: Safety to Muller



■ $\mathcal{F}_0 = \{\{0, 1, 2\}, \{0\}, \{2\}\}$

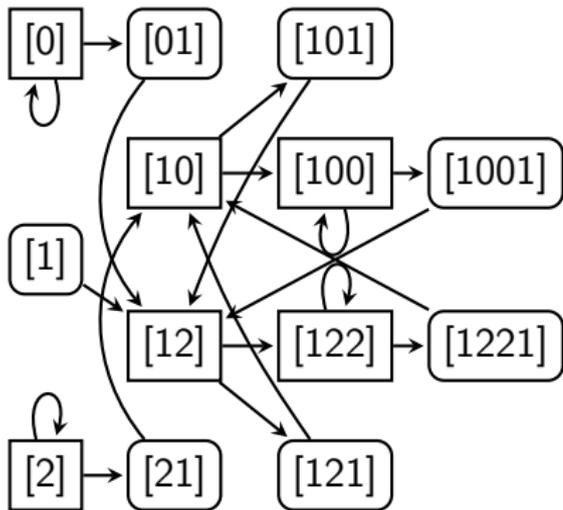
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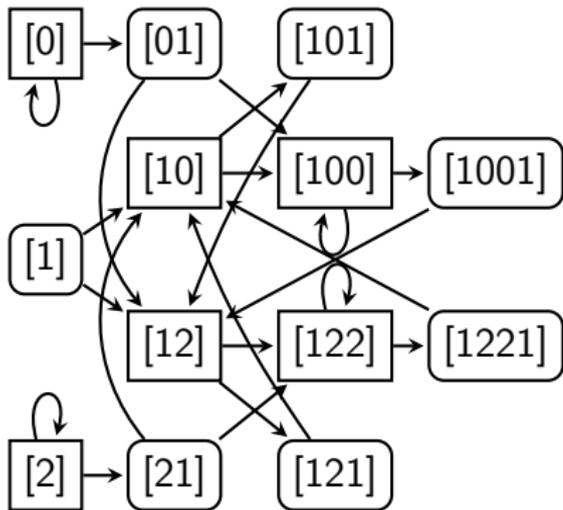


Use the winning region of safety game as memory structure..

The Proof: Safety to Muller



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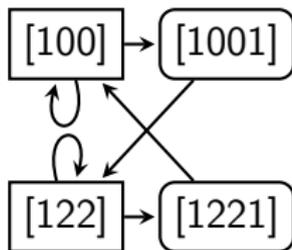


..or construct a permissive strategy..

The Proof: Safety to Muller



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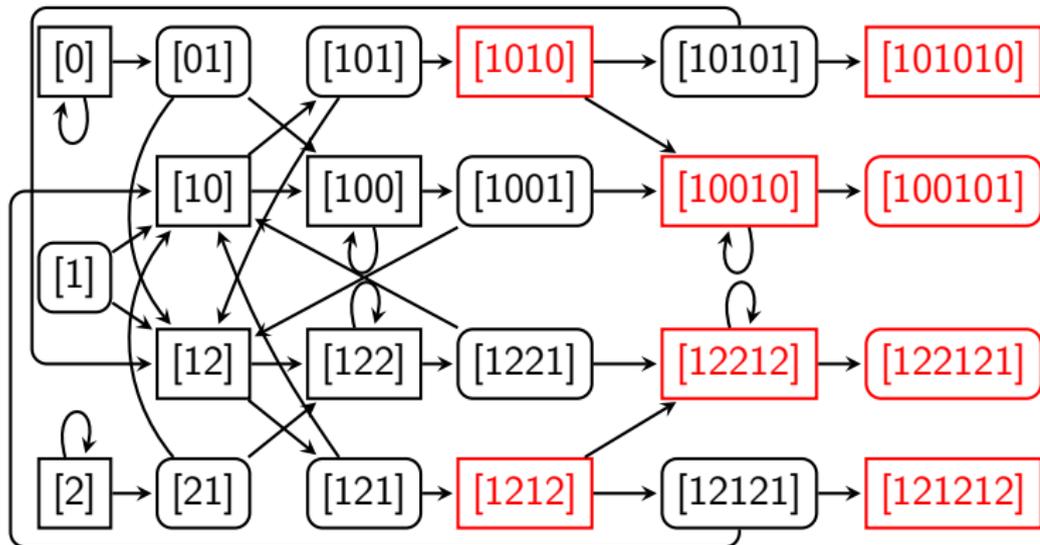
..or keep only maximal elements.

The Proof: Muller to Safety



- $\mathcal{F}_0 = \{\{0, 1, 2\}, \{0\}, \{2\}\}$

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Mimic strategy that prevents Player 1 from reaching a score of 3.

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Conclusion

Solving Muller games via safety games:

- New algorithm for Muller games: just solve the safety game.
- New memory structure for Muller games: maximal elements of winning region (antichain).
- New concept: permissive strategies for Muller games.
- Same constructions applicable for many other types of games.

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- Same constructions applicable for many other types of games.

Ongoing and future work:

- A progress measure algorithm for Muller games?
- Is there a tradeoff between size and quality of a strategy?
- Can you play infinite games in infinite arenas in finite time?