
Playing Muller Games in a Hurry

Joint work with John Fearnley, University of Warwick

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Motivation

Robert McNaughton: *Playing Infinite Games in Finite Time*. In: *A Half-Century of Automata Theory*, World Scientific (2000).

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“Winning regions should be equal”

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A Muller game $(G, \mathcal{F}_0, \mathcal{F}_1)$ consists of an arena $G = (V, V_0, V_1, E)$ and a partition $(\mathcal{F}_0, \mathcal{F}_1)$ of 2^V .

Rules:

- Players move a token through the arena ad infinitum.
- Player i wins play (infinite path) iff the set of vertices visited infinitely often is in \mathcal{F}_i .

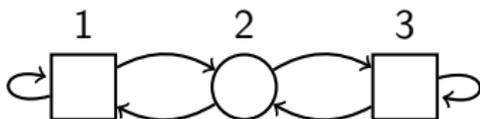
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- $\mathcal{F}_1 = 2^{\{1,2,3\}} \setminus \mathcal{F}_0$

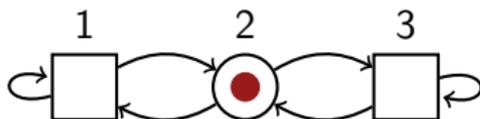
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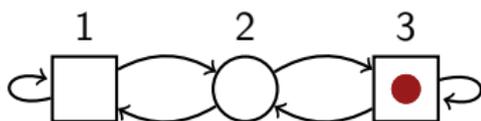
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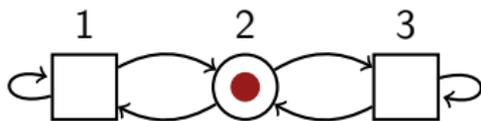
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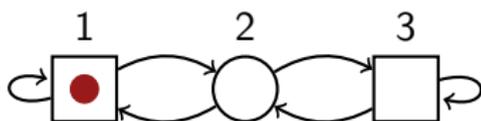
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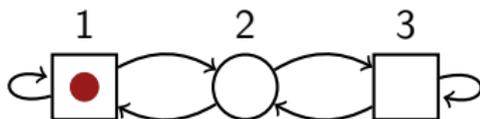
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Winning strategy for Player 0 (circles): coming from 1 to 2 move to 3, coming from 3 to 2 move to 1.

Scoring Functions

For $F \subseteq V$ define $\text{Sc}_F: V^+ \rightarrow \mathbb{N}$:

$$\text{Sc}_F(w) = \max\{k \mid \text{exist words } x_1, \dots, x_k \in V^+ \text{ s.t.} \\ x_1 \cdots x_k \text{ is suffix of } w \text{ and } \text{Occ}(x_i) = F \text{ for all } i\}$$

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For $\mathcal{F} \subseteq 2^V$ define $\text{MaxSc}_{\mathcal{F}}: V^+ \cup V^\omega \rightarrow \mathbb{N} \cup \{\infty\}$:

$$\text{MaxSc}_{\mathcal{F}}(\rho) = \max_{F \in \mathcal{F}} \max_{w \sqsubseteq \rho} \text{Sc}_F(w)$$

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$\mathcal{F} = \{\{a, b\}, \{a, b, c\}\}$:

$$\text{MaxSc}_{\mathcal{F}}(w) = 3$$

Finite-time Muller Games

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Definition

Finite-time Muller game: $(G, \mathcal{F}_0, \mathcal{F}_1, k)$ with threshold $k \geq 2$.

Rules:

- Players move a token through the arena.
- Stop play w as soon as score of k is reached for the first time.
- There is a unique F such that $\text{Sc}_F(w) = k$ (see above).
- Player i wins w iff $F \in \mathcal{F}_i$.

Results

McNaughton's version: stop play when some S_{c_F} reaches $|F|! + 1$.

Theorem (McNaughton 2000)

The winning regions in a Muller game and in McNaughton's finite-time Muller game coincide.

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Stronger statement, which implies the theorem:

Lemma

On his winning region, Player i can prevent her opponent from ever reaching a score of 3 for every set $F \in \mathcal{F}_{1-i}$.

Conclusion

Results:

	Reduction	McNaughton	here
Threshold	–	$ F ! + 1$	3
Play Length	$\leq n \cdot n! + 1$	$\leq (n! + 1)^n$	$\leq 3^n$
Space	$\mathcal{O}(n!)$	$\mathcal{O}((n! + 1)^n)$	$\mathcal{O}(3^n)$

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Open Questions:

- Is the finite-time Muller game with threshold 2 equivalent to the original Muller game?
- Given a winning strategy for a finite-time Muller game, can we turn it into a winning strategy for the Muller game?